

Algorithms for GIS

Today

- More C/compiling/Makefiles
 - Working with multiple files
 - Pointer exercises
- Announcements:
 - Assignment 1 due on Tuesday
 - Assignment 2 posted, due on 9/22
 - Josh office hours: Mon 8pm
 - Need your feedback!

Header files

- Example: implement a linked list
 - **list.h:**
 - is the interface to the outside world
 - contains type definitions and signature of functions that are meant to be used by other modules
 - **list.c:**
 - implements all functions in list.h

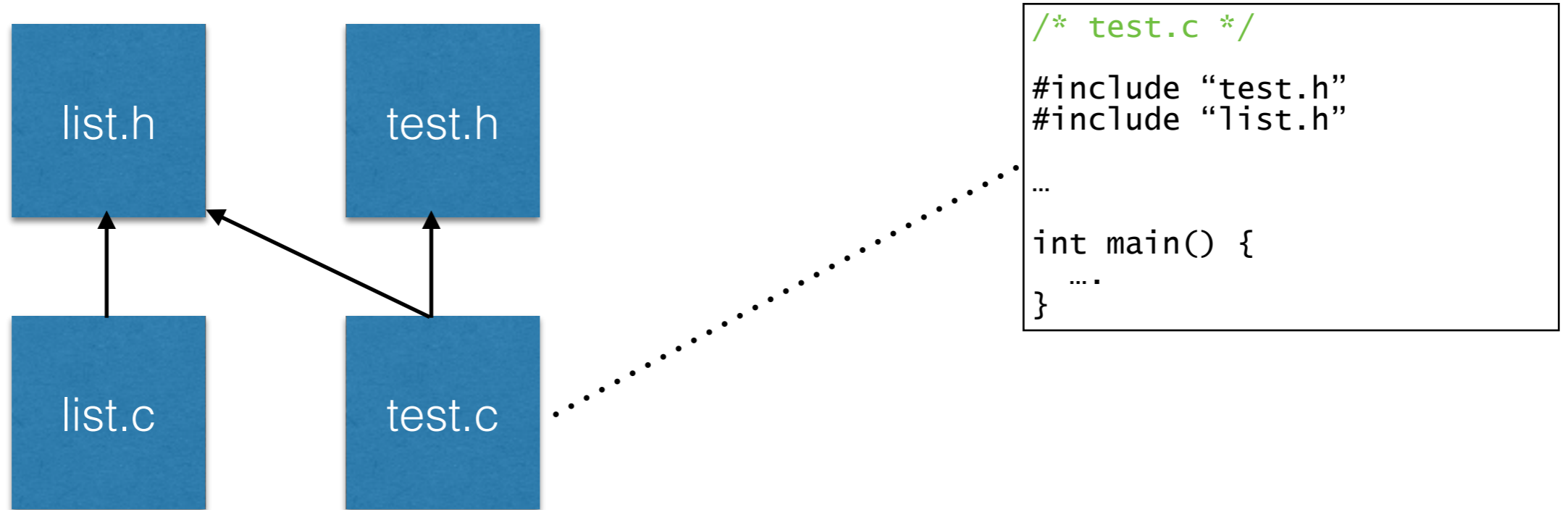
```
/* list.h */  
  
typedef struct node_t {  
    int data;  
    struct node_t* next;  
}  
  
typedef struct list_t {  
    Node* head;  
} List;  
  
List* init();  
...
```

list.h

list.c

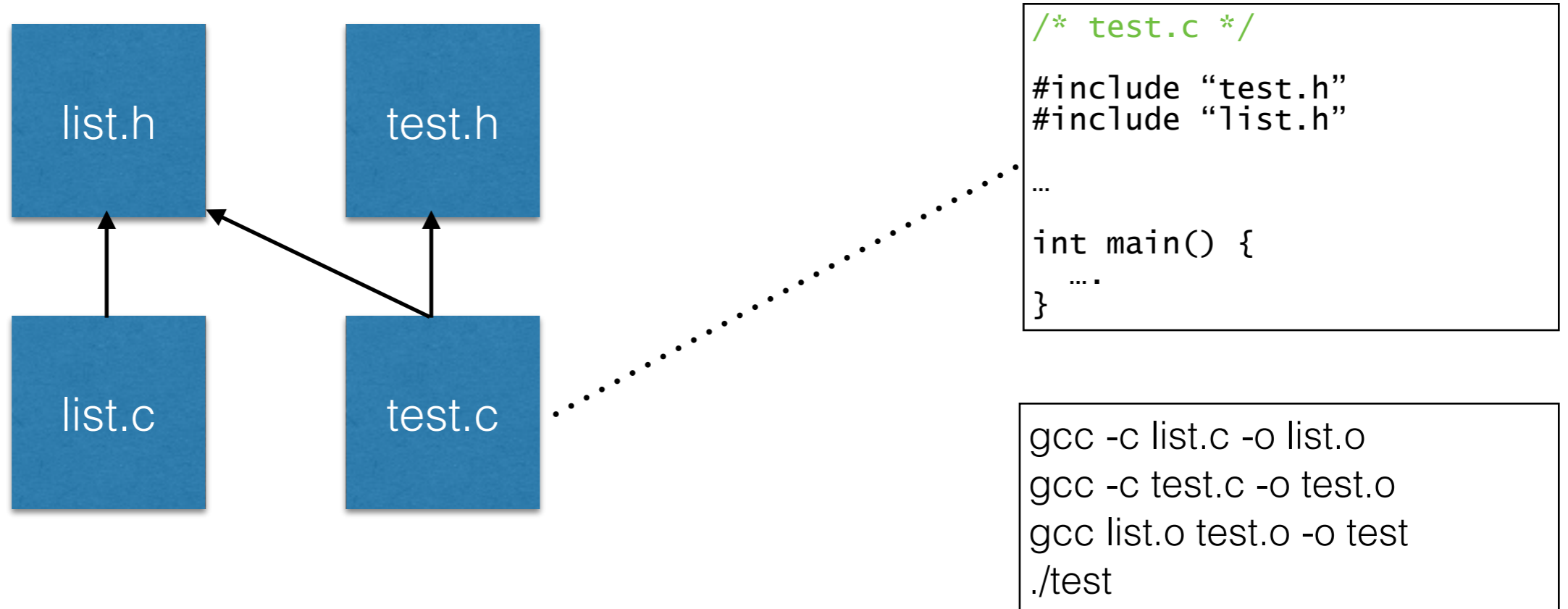
```
/* list.c */  
  
#include "list.h"  
  
List* init() {  
    //implement init  
    ...  
}
```

Working with multiple files



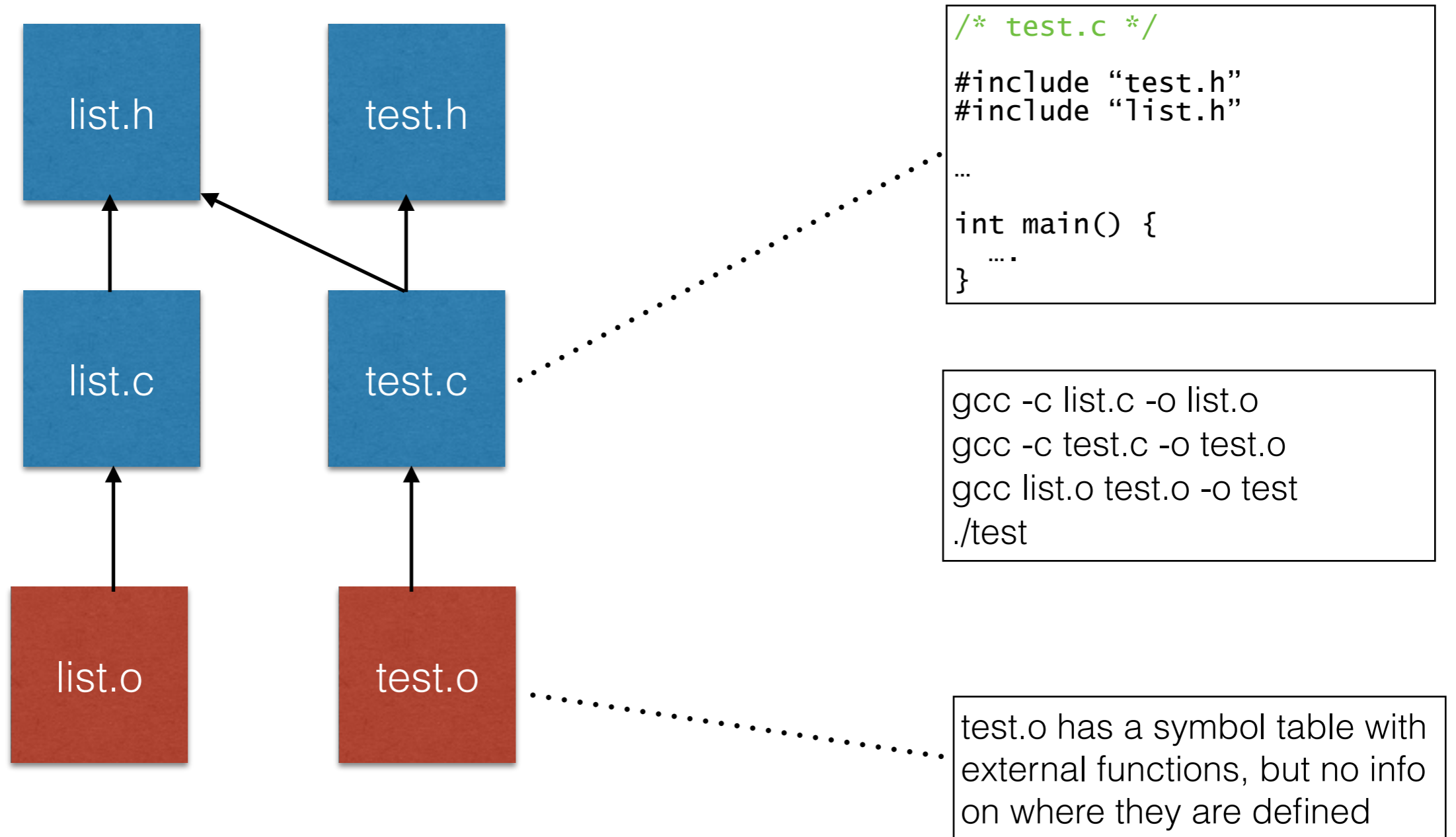
- If test.c needs to use some list functions
 - needs to `#include "list.h"`

Working with multiple files

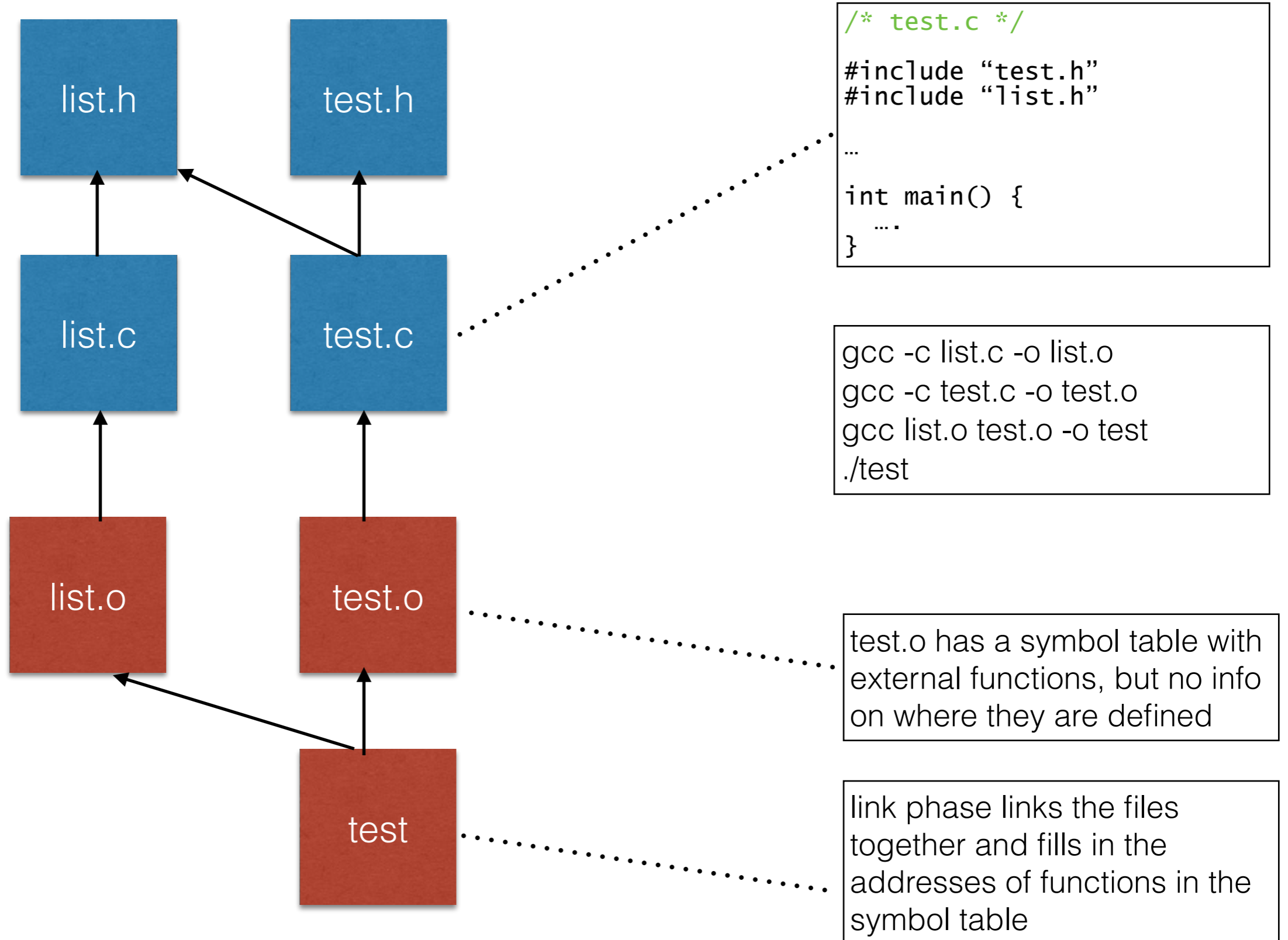


- **Compilation has 2 phases**
 - compile only (gcc -c): each xxx.c file ==> xxx.o file
 - for each file that contains a main():
 - link the .o files of the headers that it needs to create the executable

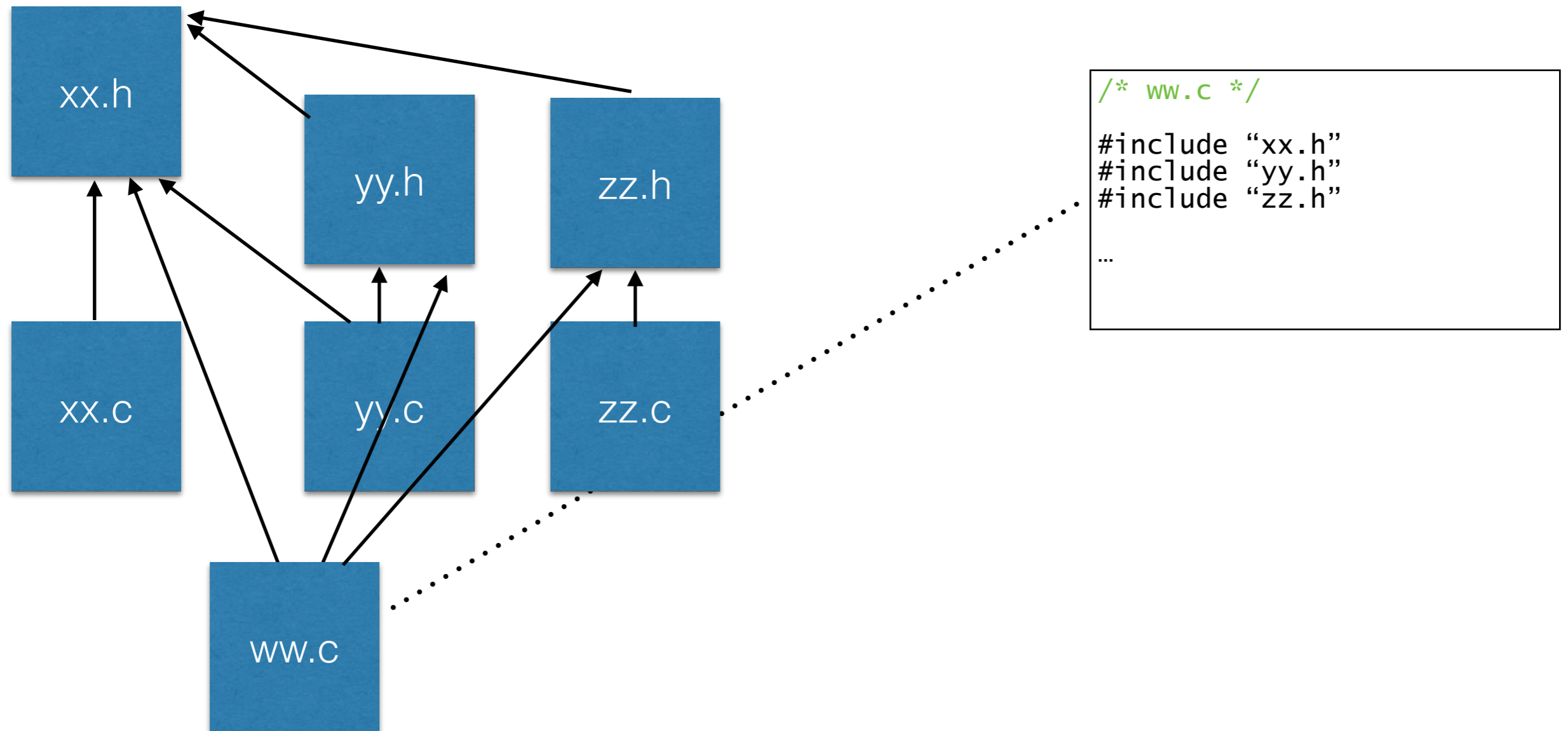
Working with multiple files



Working with multiple files

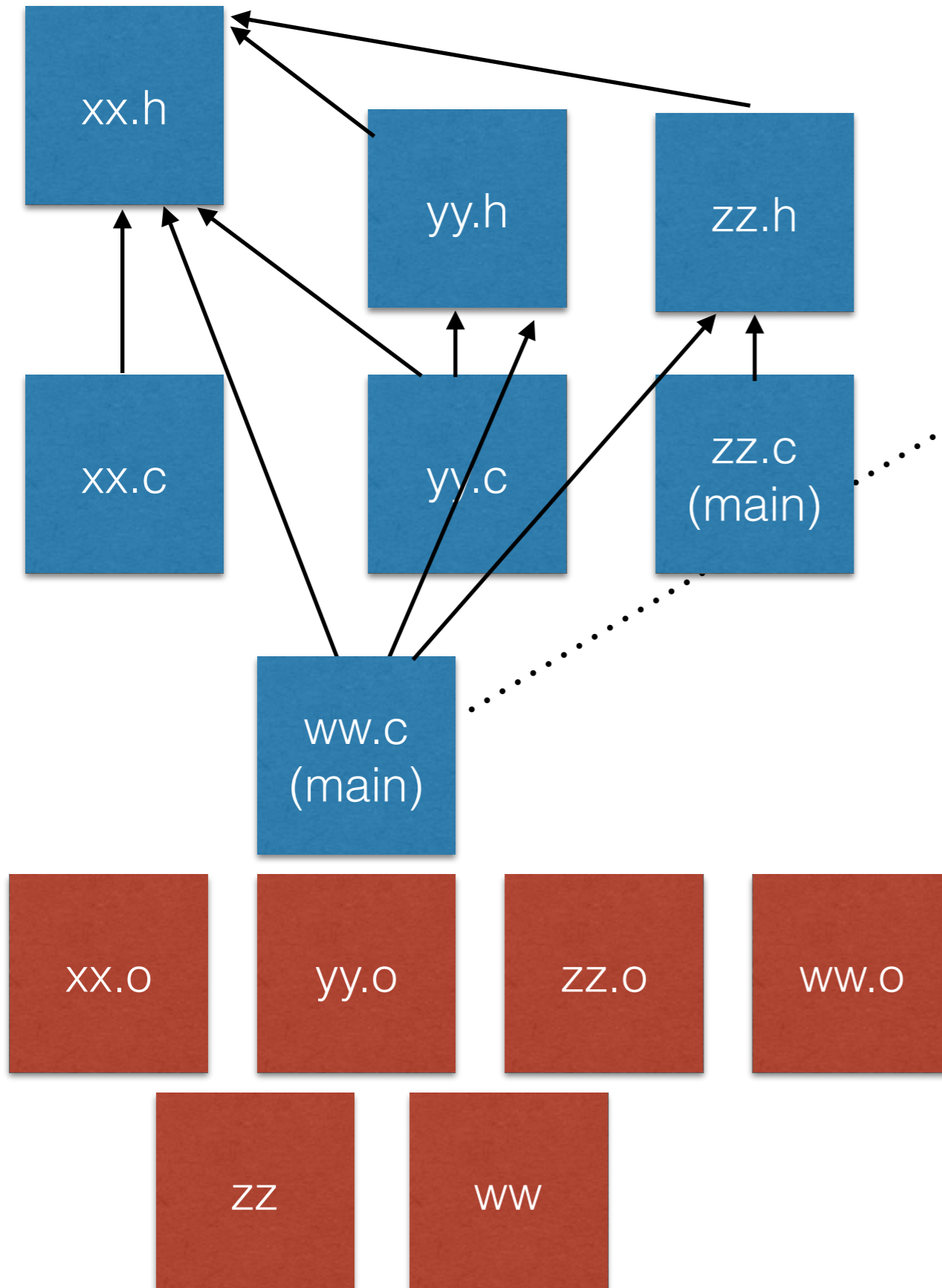


Complex dependency graph



- Each file must include all headers it needs
- The dependency graph must be ACYCLIC!
 - If not, very weird compile errors

Complex dependency graph



```
/* ww.c */  
#include "xx.h"  
#include "yy.h"  
#include "zz.h"  
...  
...
```

Why?

- For efficiency
 - compiling large projects is slow
 - if change one line in a file, you re-compile only the object files and executables that depend on it, directly or indirectly
- make utility
 - Makefile specifies dependencies
 - 'make' keeps track of when files were last modified ==> figures out what changed and what needs to be recompiled

Pointers

T x;

- Any variable is stored somewhere in memory and thus has an address, which can be retrieved with operator &

&x gives the address of variable x

- An address is called a pointer
- The address of a variable of type T is considered to have type T*

&x has type T*;

- Given an address, we might want to know what is stored at that address. That's called dereferencing the pointer, and it's done with operator *

T* p;

*p is of type T, and it is the value stored at address p

- **Caveat: Dereferencing an invalid address is a **BUG**.**
- A bug of this type doesn't always manifest, and does not manifest in the same way. That is, it might give you a segfault. Or not. Still, your program has a bug in it and its behavior is unpredictable.

Pointers

- Rule: make sure you initialize a pointer before you dereference it
 - by assigning it the address of a variable
 - by calling malloc()
 - by assigning it the value of another valid pointer

Pointers

- Rule: make sure you initialize a pointer before you dereference it
 - by assigning it the address of a variable
 - by calling malloc()
 - by assigning it the value of another valid pointer
- Perhaps this is boring.. Consider this.
 - You WILL get segfaults
 - You will spend a LONG time figuring it out
 - It's ALWAYS because you break this one rule
- And remember
 - Bad memory references do not always manifest
 - The program might work fine on one computer, but not on other.

Pointer quiz

- We want to write a function to allocate an array of n element of type T
- We'll write it two ways:
 - return the array
 - take the array as parameter

Is this working?

```
//assume T is a type

T* create(int n) {
    T* result;
    result = (T*)malloc(n*sizeof(T));
    assert(result);
    return result;
}

int n=100;
T* x;
x = create(n);
//is x an array of 100 elements?
```

```
//assume T is a type

void create(int n, T* a) {
    a = (T*) malloc(n*sizeof(T));
    assert(a);
}

int n=100;
T* x;
create(n, x);
//is x an array of 100 elements?
```