

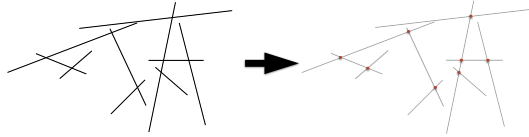
Line segment intersection

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1

Line segment intersection

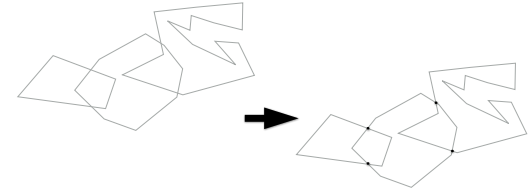
Problem: Given a set of line segments in 2D, find all their pairwise intersections.



2

Line segment intersection

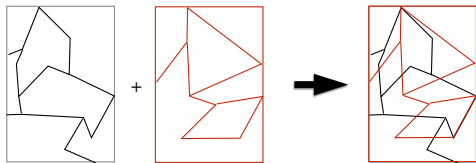
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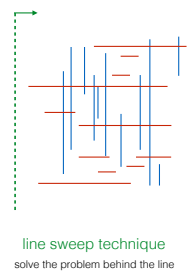
Line segment intersection

Problem: Given a set of line segments in 2D, find all their pairwise intersections.

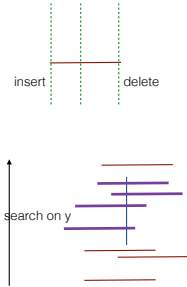


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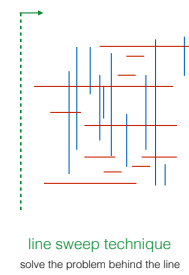
Orthogonal line segment intersection



5



Orthogonal line segment intersection



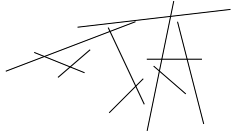
6

- n : size of the input (number of segments)
- k : size of output (number of intersections)

Result: The intersections of a set of n orthogonal segments in the plane can be found in $O(n \lg n + k)$ time.

General line segment intersection

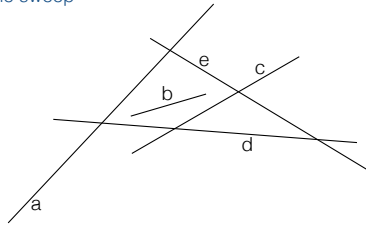
- n : size of the input (number of segments)
- k : size of output (number of intersections)



- Extend sweep line idea
- We'll get an overall bound of $O(n \lg n + k \lg n)$ which improves on the naive $O(n^2)$ when k is small
- The algorithm was developed by Jon Bentley and Thomas Ottmann in 1979
- Simple (once you see it) and practical

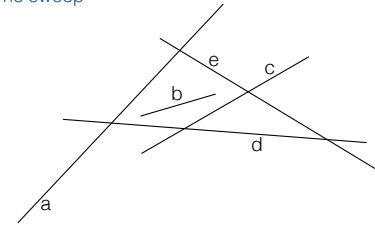
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The sweep



8

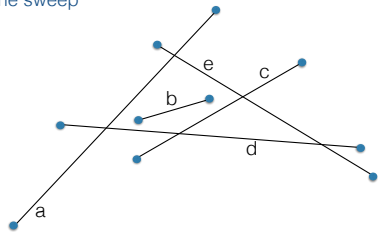
The sweep



- Let X be the set of all x-coords of segments

9

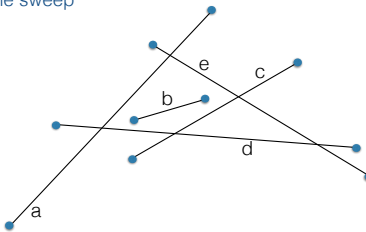
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- Let X be the set of all x-coords of segments

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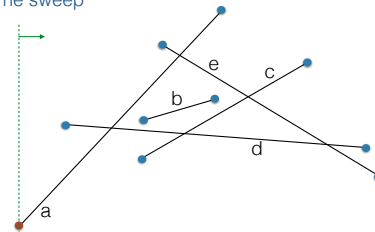
The sweep



- Let X be the set of all x-coords of segments
- Traverse the events in X in order

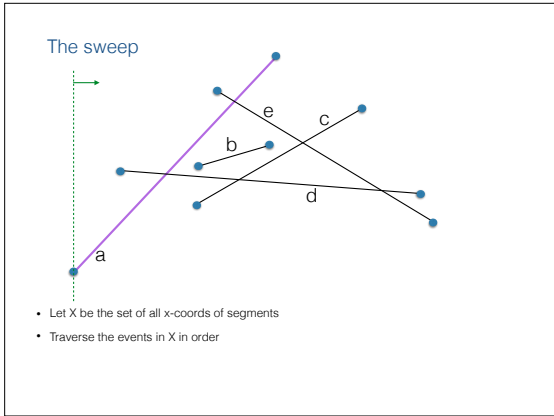
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The sweep

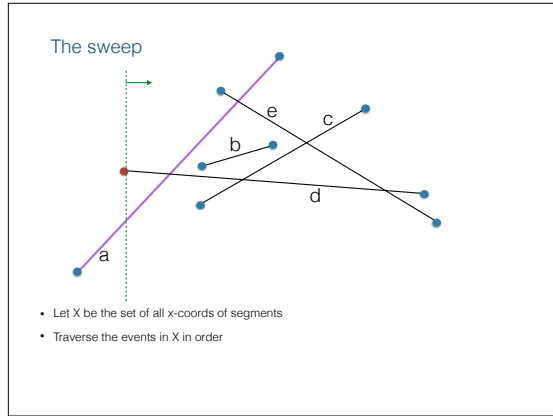


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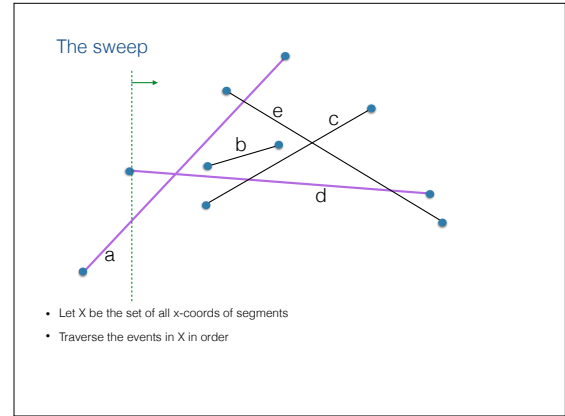
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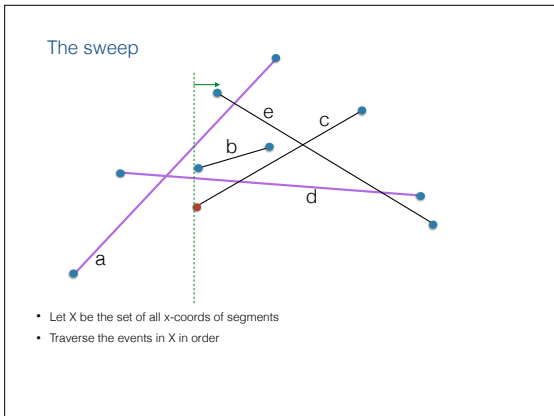
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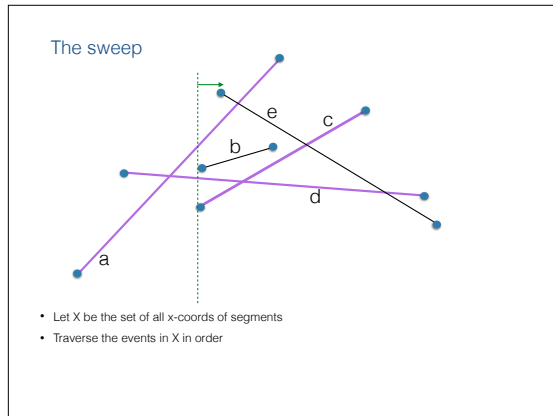
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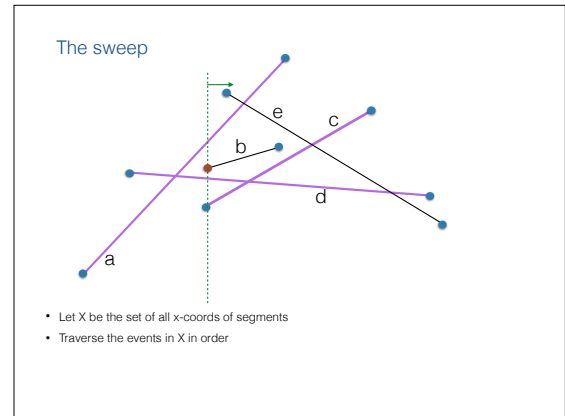
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16



17



18

The sweep

- Let X be the set of all x-coords of segments
- Traverse the events in X in order

19

The sweep

- 4 segments are active
- How do we order these segments?
- How do we detect intersections?

20

How do we order these segments?

a below d a above d

general segments orthogonal segments

21

Key idea #1

- How do we order the active segments?
 - Use above-below order
 - Order will flip at intersection point

22

Key idea #2

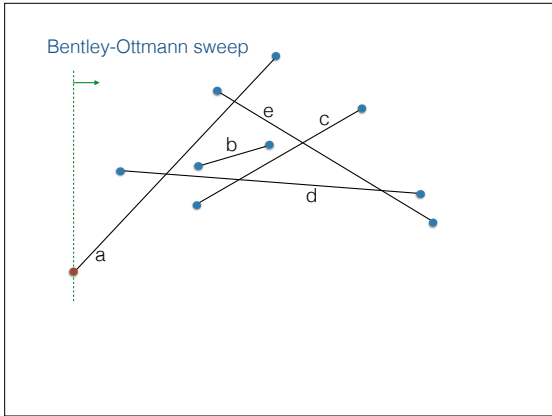
- How do we detect intersections?
 - Segments that intersect are consecutive in above-below order just before they intersect

23

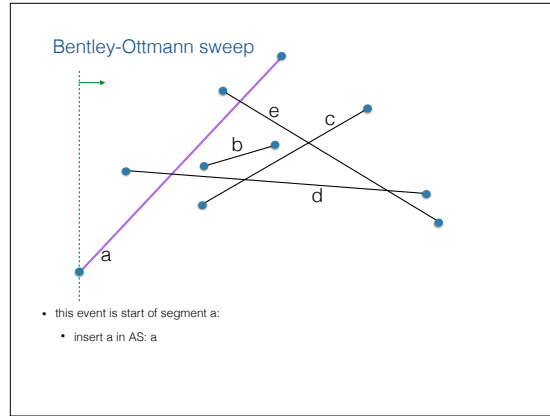
Bentley-Ottmann sweep

- Let X be the set of all x-coords of segments
- Initialize $AS = \{\}$
- Traverse events in order

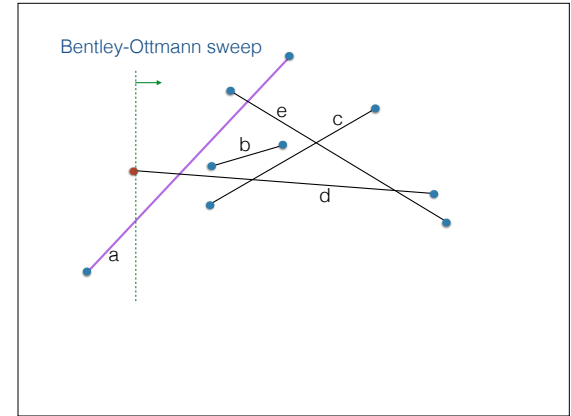
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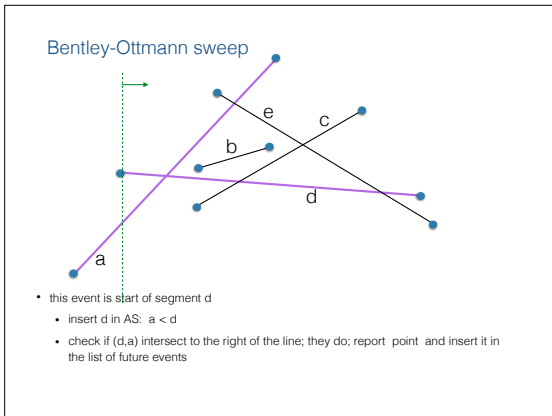
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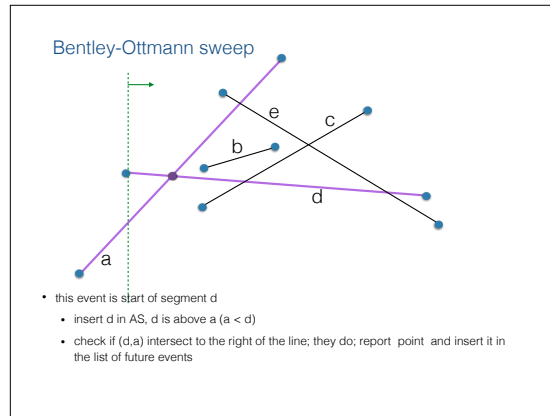
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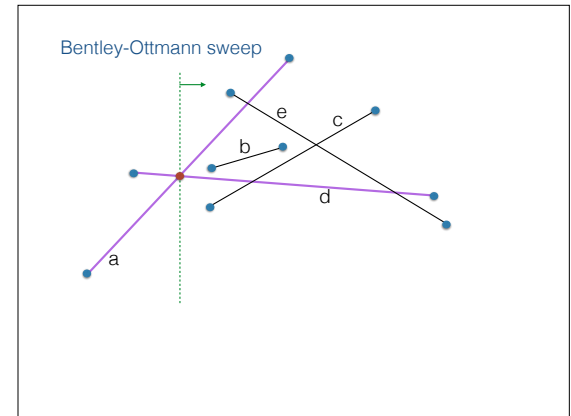
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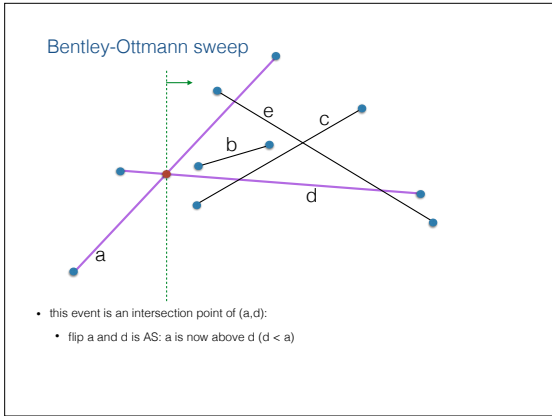
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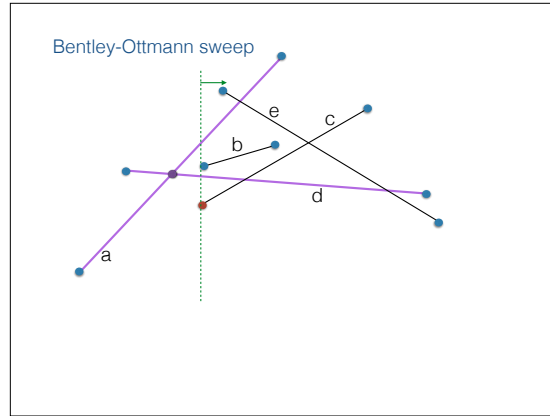
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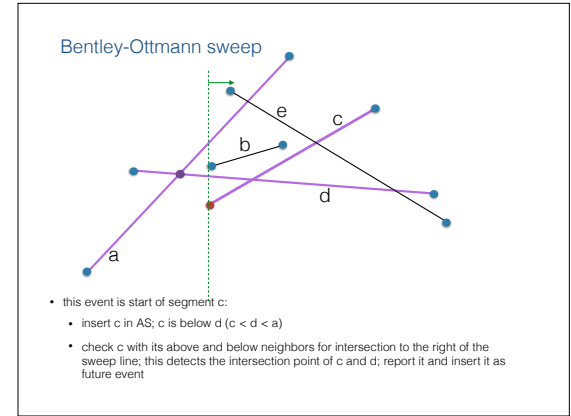
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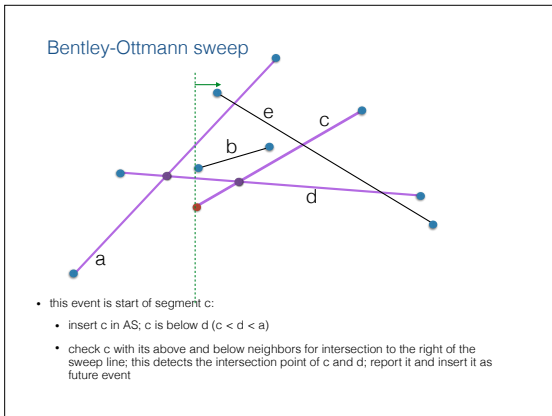
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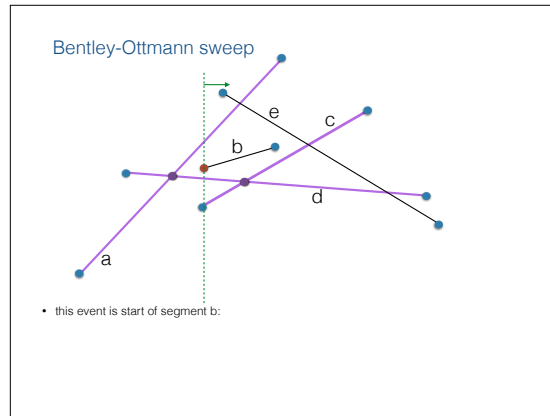
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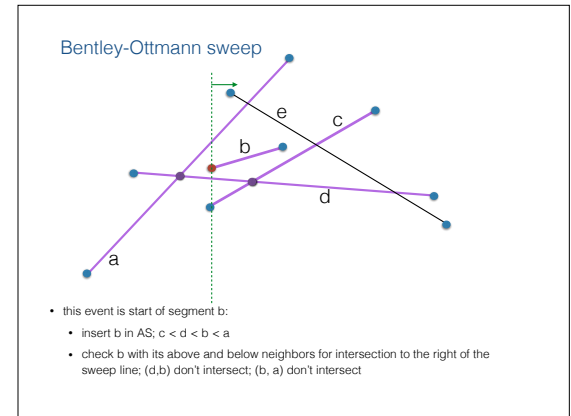
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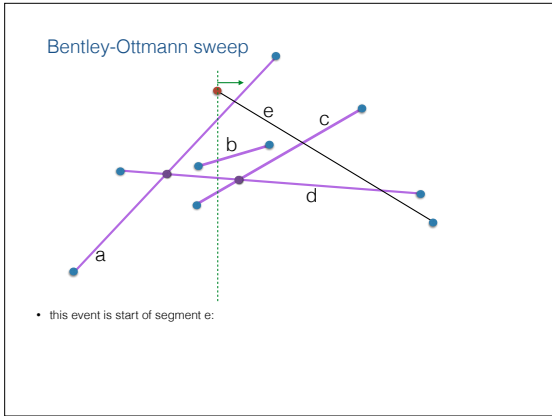
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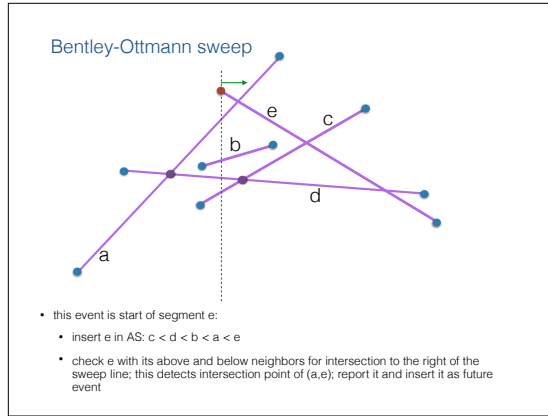
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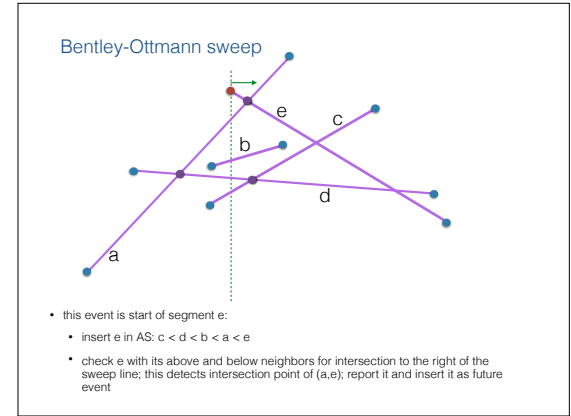
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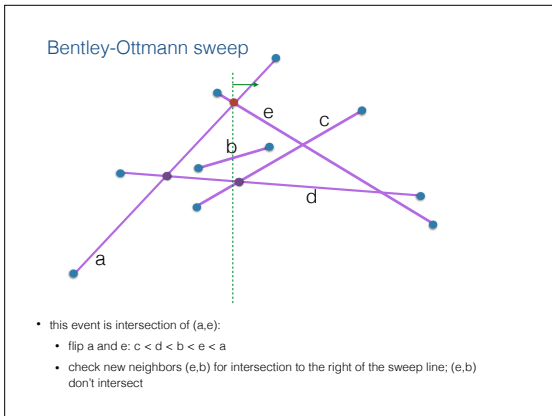
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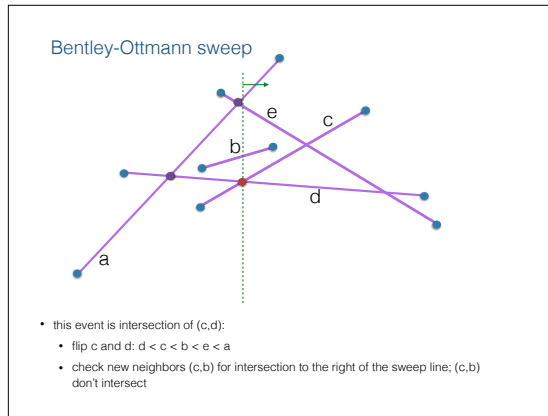
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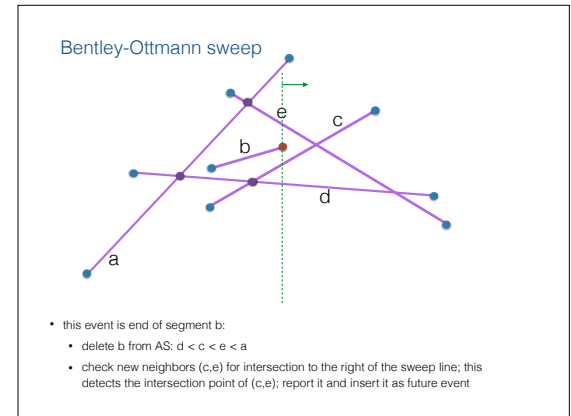
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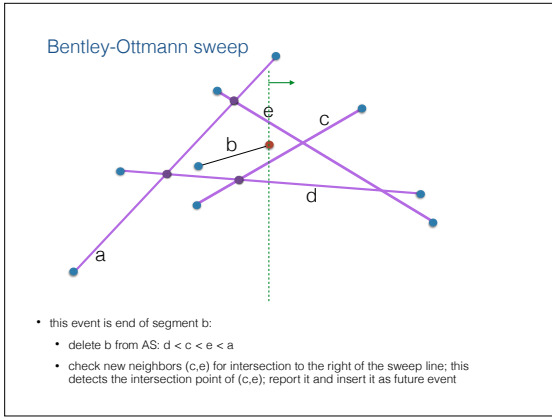
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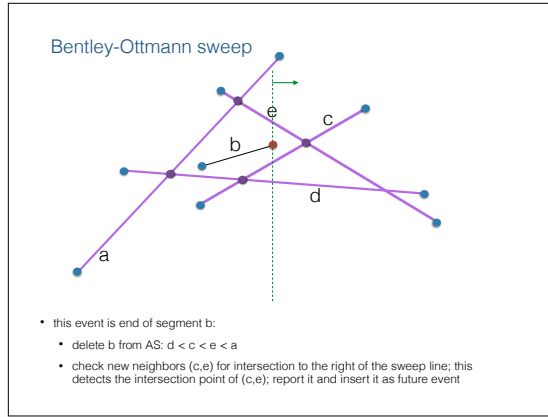
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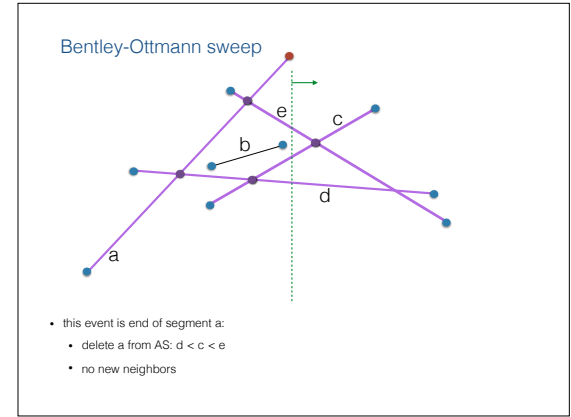
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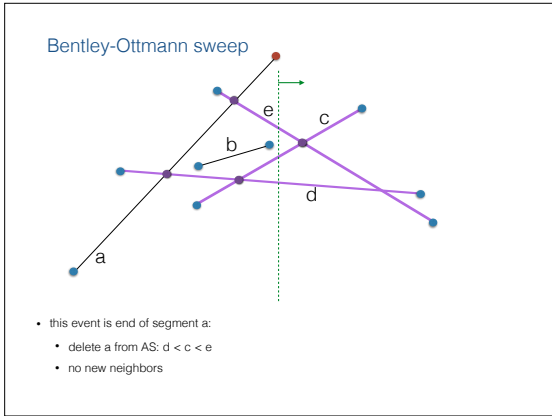
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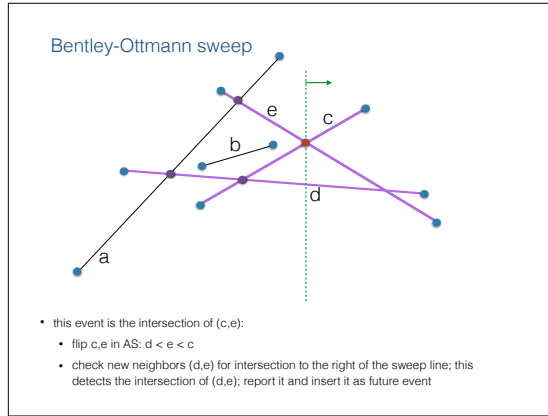
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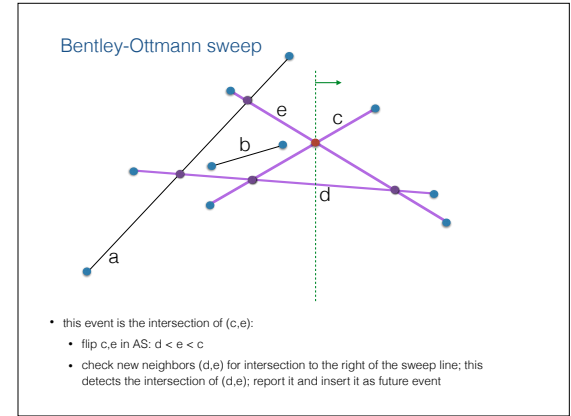
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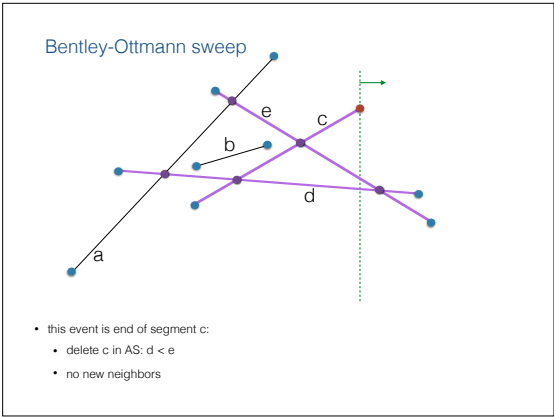
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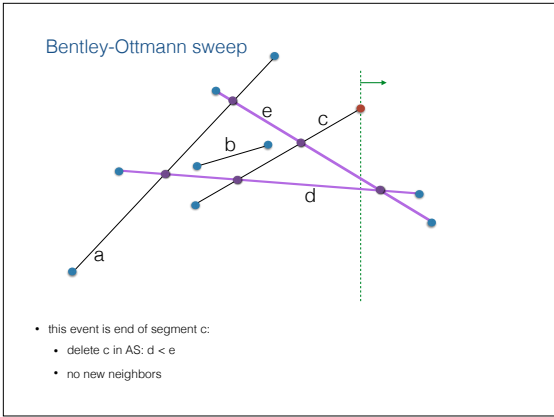
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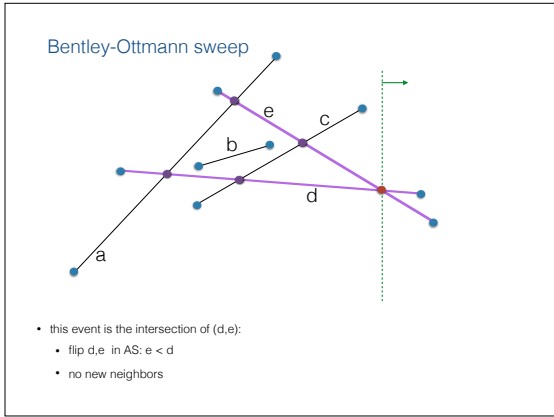
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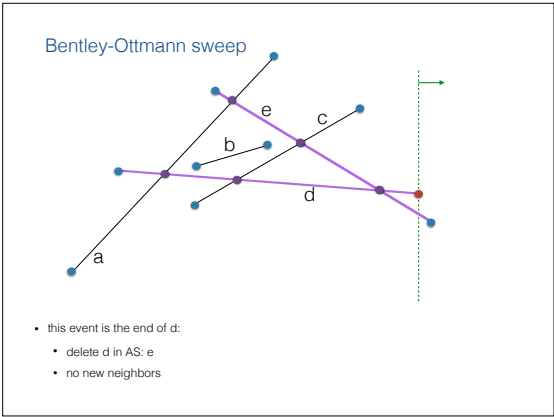
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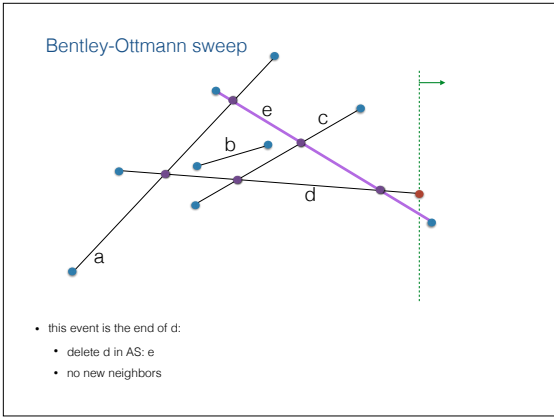
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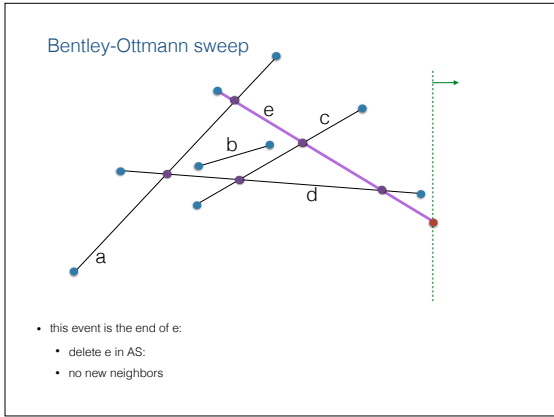
51



52



53



54

Bentley-Ottmann sweep

- this event is the end of e:
 - delete e in AS:
 - no new neighbors

55



56

Bentley-Ottmann sweep

- Simplifying assumptions
 - no vertical segments
 - no two segments intersect at their endpoints
 - no three (or more) segments have a common intersection
 - all endpoints (of segments) and all intersection points have different x-coordinates
 - no segments overlap
- These assumptions are not realistic for real data.
- But, they don't provide insight into the plane sweep technique, so we omit them
- Real data challenges
 - dealing with degenerate cases
 - dealing with finite precision arithmetic and precision problems

57

Bentley-Ottmann sweep

We'll maintain the following invariants during the algorithm:

58

Bentley-Ottmann sweep

We'll maintain the following invariants during the algorithm:

- Active structure AS:
 - For any position of the sweep line SL, AS contains all active segments (ie segments that start before SL and end after SL)
 - AS is sorted by their y-coordinates of their intersection with SL

59

Bentley-Ottmann sweep

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- Active structure AS:
 - For any position of the sweep line SL, AS contains all active segments (ie segments that start before SL and end after SL)
 - AS is sorted by their y-coordinates of their intersection with SL
- Event list EL:
 - For any position of SL, EL contains segment endpoints to the right of SL, and also the intersections to the right of SL of active segments that were/are neighbors in SL
 - EL is sorted by x-coordinate

60

Bentley-Ottmann sweep

We'll maintain the following invariants during the algorithm:

- Active structure AS:
 - For any position of the sweep line SL, AS contains all active segments (ie segments that start before SL and end after SL)
 - AS is sorted by their y-coordinates of their intersection with SL
- Event list EL:
 - For any position of SL, EL contains segment endpoints to the right of SL, and also the intersections to the right of SL of active segments that were/are neighbors in SL
 - EL is sorted by x-coordinate
- For any position of the sweep line SL, all pairs of intersecting dead segments have been reported.

61

Algorithm Bentley-Ottmann (S)

```
//S is a set of n line segments in the plane
• initialize AS= {}
• sort 2n endpoints of all segments in S by x-coord and store them in EventList
• while EventList not empty
  • let e be the next event from EventList; delete it from EL
  //sweep line moves to x=e.x
  • if e is left endpoint of a segment l
    // becomes active
    • insert l in AS in the right place
    • check if l intersects with l->prev and l->succ in AS to the right of the sweep line; if they do, insert their intersection point in the EventList
    //optional: since l.prev and l.succ are not neighbors anymore, we check if they intersect and if they do, delete that intersection point from the EventList
  • if e is the right endpoint of a segment
    • ...
  • if e is the intersection of two segments
    • ...
  • end.
```

62

Bentley-Ottmann sweep

Questions

- AS
 - What operations do we do on AS?
 - What data structure should we use for AS?
- EL
 - Note that we know a priori the 2n events corresponding to start and end-points of segments, but EL is not static; the events corresponding to intersection points are generated on the fly
 - What operations do we do on EL?
 - What data structure should we use for EL?

63

Bentley-Ottmann sweep

Running time

- AS
 - What is the size of AS?
 - $O(n)$
 - How many operations?
 - $O(n+k)$
 - Overall time?
 - $O((n+k)\lg n)$
- EventList
 - What is the size of EventList?
 - $O(n+k)$
 - How many operations?
 - $O(n+k)$
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64

Bentley-Ottmann sweep

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Bentley-Ottmann sweep runs in $O((n+k)\lg n)$ time.

65