Computational Geometry [csci 3250]

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What does the guard see?



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We say that a set of guards **covers** polygon P if every point in P is visible to at least one guard.



Does the point guard the triangle?

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Can any triangle be guarded with one point?

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Does the point guard the quadrilateral?

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Can any quadrilateral be guarded with one point?



Questions:

1. Given a polygon P of size n, what is the smallest number of guards (and their locations) to cover P?



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- 1. Given a polygon P of size n, what is the smallest number of guards (and their locations) to cover P? NP-Complete
- 2. **Klee's problem:** Consider all polygons of n vertices, and for each one, the smallest number of guards to cover it. What is the worst-case?

Notation

'n-gon"

- Let Pn: polygon of n vertices
- Let g(P) = the smallest number of guards to cover P
- Let $G(n) = max \{ g(P_n) | all P_n \}.$
- What does this mean?
 - G(n) is the smallest number that always works for any n-gon. It is sometimes necessary and always sufficient to guard a polygon of n vertices.
 - G(n) is necessary: there exists a Pn that requires G(n) guards
 - G(n) is sufficient: any P_n can be guarded with G(n) guards
- Klee's problem: find G(n)

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Our goal (i.e. Klee's goal) is to find G(n).

Trivial bounds

- G(n) >= 1: obviously, you need at least one guard.
- G(n) <= n : place one guard in each vertex

n=3



Any triangle needs at least one guard. One guard is always sufficient.

$$G(3) = 1$$

n=4



Any quadrilateral needs at least one guard. One guard is always sufficient.

G(4) = 1



Can all 5-gons be guarded by one point?



G(5) = 1



n=6





n=6



G(6) = 2

G(n) = ?

Come up with a P_n that requires as many guards as possible.



How many guards does this need?



This polygon requires $\lfloor n/3 \rfloor$ guards => G(n) >= $\lfloor n/3 \rfloor$



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Are there P_n that require more guards, or, are [n/3]guards always sufficient for any P_n?

It was shown that $\lfloor n/3 \rfloor$ is always sufficient for any $P_{n:}$

Any P_n can be guarded with at most $\lfloor n/3 \rfloor$ guards.

- (Complex) proof by induction
- Subsequently, simple and beautiful proof due to Steve Fisk, who was Bowdoin Math faculty.
- Proof in The Book.

Proofs from THE BOOK

From Wikipedia, the free encyclopedia

Proofs from THE BOOK is a book of mathematical proofs by Martin Aigner and Günter M. Ziegler. The book is dedicated to the mathematician Paul Erdős, who often referred to "The Book" in which God keeps the most elegant proof of each mathematical theorem. During a lecture in 1985, Erdős said, "You don't have to believe in God, but you should believe in The Book."

Content [edit]

Proofs from THE BOOK contains 32 sections (44 in the fifth edition), each devoted to one theorem but often containing multiple proofs and related results. It spans a broad range of mathematical fields: number theory, geometry, analysis, combinatorics and graph theory. Erdős himself made many suggestions for the book, but died before its publication. The book is illustrated by Karl Heinrich Hofmann. It has gone through five editions in English, and has been translated into Persian, French, German, Hungarian, Italian, Japanese, Chinese, Polish, Portuguese, Korean, Turkish, Russian and Spanish.

The proofs include:

- Proof of Bertrand's postulate
- · Proof that e is irrational (also showing the irrationality of certain related numbers)
- Six proofs of the infinitude of the primes, including Euclid's and Furstenberg's
- · Monsky's theorem (4th edition)
- · Wetzel's problem on families of analytic functions with few distinct values
- Steve Fisk's proof of the The art gallery theorem

Deferences / m

Fisk's proof at a glance:

- 1. Any simple polygon can be triangulated.
- 2. A triangulated simple polygon can be 3-colored.
- 3. Observe that placing the guards at all the vertices assigned to one color guarantees the polygon is covered.
- 4. There must exist a color that's used at most n/3 times. Pick that color and place guards at the vertices of that color.

Given a simple polygon P, a **diagonal** is a segment between 2 nonadjacent vertices that lies entirely within the interior of the polygon.



Claim: Any simple polygon can be triangulated.

Proof idea: induction using the existence of a diagonal.



Claim 1: Any simple polygon contains at least one convex vertex





Claim 2: Any simple polygon contains at least one diagonal.



- 1. Any simple polygon can be triangulated
- 2. Any triangulation of a simple polygon can be 3-colored.



Coloring

• A coloring of a graph is an assignment of colors to vertices such that no two adjacent vertices (vertices connected by an edge) have the same color

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- A coloring of a graph is an assignment of colors to vertices such that no two adjacent vertices (vertices connected by an edge) have the same color
- The chromatic number of a graph G, $\chi(G)$
 - $\chi(G)$ = the smallest nb of colors needed to color G
- Fundamental problem in graph theory
- NP-complete to compute $\chi(G)$
- Results:
 - Any planar graph can be 5-colored. O(n) time.
 - Any planar graph can be 4-colored (proof by computer). O(n²) time.
 - Can G be 3-colored? NP-complete

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- Placing guards at vertices of one color covers P.
- Pick least frequent color! At most n/3 vertices of that color.



- 1. Any polygon can be triangulated
- 2. Any triangulation can be 3-colored
- 3. Observe that placing the guards at all the vertices assigned to one color guarantees the polygon is covered.
- 4. There must exist a color that's used at most n/3 times. Pick that color and place guards at the vertices of that color.

Claim: The set of red vertices covers the polygon. The set of blue vertices covers the polygon. The set of green vertices covers the polygon.

Because...

There are n vertices colored with one of 3 colors.

Claim: There must exist a color that's used at most n/3 times.

Proof:

Theorem: Any triangulation can be 3-colored.

Proof: