## Week 2: The closest pair of points in the plane

The problem: Given an array $P$ of $n$ points in the plane, find the closest pair. Assume that the distance between two points is given by the Euclidian distance.

## Questions

1. Formulate the 1D version of the closest pair. How can you solve it, and how fast? Try to extend this solution to the 2D problem: does it work?

For the remaining problems we consider the 2D version.
2. Describe how you can find a vertical line that splits $P$ in half. How long does this take?
3. Consider the (refined) closest pair algorithm which takes as arguments the points in $P$ sorted in two different ways. Let $P_{X}$ and $P_{Y}$ denote the points in $P$ sorted by their x- and ycoordinates, respectively. Furthermore, Let $L$ be the vertical line that splits $P$ into two halves, and let $P_{1}$ and $P_{2}$ be the set of points in $P$ to the left/right of this line, respectively.
(a) Given $P_{X}$ and $P_{Y}$, how can you find the x-coordinate of line $L$ ?
(b) Given $P_{X}$ and $P_{Y}$, how can you find $P_{1 X}$ (the points in $P_{1}$ sorted by their x-coordinates) and $P_{2 X}$ (the points in $P_{2}$ sorted by their x-coordinates)?
(c) Given $P_{X}$ and $P_{Y}$, how can you find $P_{1 Y}$ (the points in $P_{1}$ sorted by their y-coordinates) and $P_{2 Y}$ (the points in $P_{2}$ sorted by their y-coordinates)?
4. The divide-and-conquer algorithm is guaranteed to run in $O(n \lg n)$ time for any set of points in the plane. In practice it is sometimes the case that data is nice; put differently we can make certain assumptions about the data, and exploit these assumptions to come up with simpler and more efficient algorithms.
Assume that the set of points $P$ is uniformly distributed. Can you come up with a different idea to find the closest pair? Can you get $O(n)$ time?
Hint: throw a grid over the points. For the sake of the analysis, assume a grid of $k-b y-k$ cells. How many points do you expect to fall in each cell, on the average? What value of $k$ would you pick?

