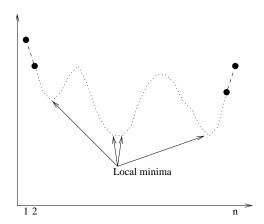
## CSci 231 Homework 4

Selection and Divide-and-conquer

## CLRS Chapter 9

- 1. (CLRS 9.3-5) Suppose that you have a "black-box" worst-case linear-time median subroutine. Give a simple, linear-time algorithm that solves the selection problem SELECT(i) for an arbitrary i.
- 2. (CLRS 9.3-7) Describe an O(n) algorithm that, given a set S of n distinct numbers and a positive integer  $k \leq n$ , determines the k numbers in S that are closest to the median of S.
- 3. Let A be a list of n (not necessarily distinct) integers. Describe an O(n)-algorithm to test whether any item occurs more than  $\lceil n/2 \rceil$  times in A. Your algorithm should use O(1) additional space.
- 4. (CLRS 9.3-6) Give an  $O(n \lg k)$  algorithm to find the k-1 elements in a set that partition the set into (approx.) k equal-sized sets  $A_1, A_2, \ldots, A_k$  such that all elements in  $A_i$  are smaller than all elements in  $A_{i+1}$ . Assume k is a power of 2.
- 5. CLRS 2-4 (The inversion problem).
- 6. Consider an array A of length n for which we know that  $A[1] \ge A[2]$  and  $A[n-1] \le A[n]$ . We say that A[x] is a *local minimum* if  $A[x-1] \ge A[x]$  and  $A[x] \le A[x+1]$ . Note that A must have at least one local minimum.



We can obviously find a local minimum in O(n) time by scanning through A. Describe an  $O(\log n)$  algorithm for finding a local minimum.