CPS 130 Homework 13 - Solutions

1. (CLRS 14.1-5) Given an element x in an n-node order-statistic tree and a natural number i, how can the ith successor of x in the linear order of the tree be determined in $O(\log n)$ time?

   Solution: The data structure should support the following two operations: OS-RANK(T, x), which returns the position of x in the linear order determined by an inorder tree walk of T in $O(\lg n)$ time, and OS-SELECT(x, i), which returns a pointer to the node containing the ith smallest key in the subtree rooted at x in $O(\lg n)$ time. The ith successor of x is given by OS-SELECT(X,OS-RANK(T, x)+I), which will also run in $O(\lg n)$ time.

2. (CLRS 14.2-1) Show how the dynamic-set queries MINIMUM, MAXIMUM, SUCCESSOR and PREDECESSOR can each be supported in $O(1)$ worst-case time on an augmented order-statistic tree. The asymptotic performance of other operations should not be affected. (*Hint: Add pointers to nodes.*).

   Solution:

3. In this problem we consider a data structure for maintaining a multi-set M. We want to support the following operations:

   - $Init(M)$: create an empty data structure M.
   - $Insert(M, i)$: insert (one copy of) i in M.
   - $Remove(M, i)$: remove (one copy of) i from M.
   - $Frequency(M, i)$: return the number of copies of i in M.
   - $Select(M, k)$: return the k'th element in the sorted order of elements in M.

   If for example M consists of the elements
   $$< 0, 3, 3, 4, 4, 7, 8, 8, 8, 9, 11, 11, 11, 11, 13 >$$
   then $Frequency(M, 4)$ will return 2 and $Select(M, 6)$ will return 7.

   Let $|M|$ and $\|M\|$ denote the number of elements and the number of *different* elements in M, respectively.

   (a) Describe an implementation of the data structure such that $Init(M)$ takes $O(1)$ time and all other operations take $O(\log \|M\|)$ time.

   Solution: The idea is to store the **distinct** elements of the multi-set in ared-black tree. For each node x in the tree which stores the value k maintain a counter $c(x) =$ how many elements in the multi-set are equal to k. $Init(M)$ simply initializes the red-black tree. $Insert(M, i)$ first searches for i in the tree: if it exists, it increments its counter, otherwise it inserts it and sets its counter to 1. $Remove(M, i)$ searches for i in the tree and if it exists, it decrements its counter, and if the counter becomes

0 it deletes that node from the tree. $Frequency(M, i)$ searches for i and returns its counter.

In order to implement $Select(M, k)$ we need to augment the tree with extra information such that each node can find out its rank. This is basically the same problem as augmenting a red-black tree in order to answer order statistics queries in $O(\lg n)$ time. We store in each node x a field $size(x)$ which is the total number of nodes in the subtree rooted at x, which can be computed as

$$size(x) = size(left(x)) + size(right(x)) + counter(x).$$

(b) Design an algorithm for sorting a list L in $O(|L| \log \|L\|)$ time using this data structure.

Solution: Insert each element from the list into this data structure and then select each element.

$$\texttt{For } i = 1 \texttt{ to } |L| \;\; Insert(M, L[i])$$
$$\texttt{For } i = 1 \texttt{ to } |L| \;\; Select(M, i)\texttt{.}$$

As the tree will contain $\|L\|$ distinct elements, each call of $Insert()$ or $Select()$ will take $O(\log \|L\|)$ time.