Lecture 9: Heaps. Heapsort.

(CLRS 6)

May 28th, 2002

1 Introduction

- We have discussed several fundamental algorithms (e.g. sorting)
- We will now turn to *data structures*; Play an important role in algorithms design.
  - Today we will discuss priority queues and next time structures for maintaining ordered sets.

2 Priority Queue

- A priority queue supports the following operations on a set S of n elements:
  - INSERT: Insert a new element e in S
  - FINDMIN: Return the minimal element in S
  - DELETEMIN: Delete the minimal element in S
- Sometimes we are also interested in supporting the following operations:
  - CHANGE: Change the key (priority) of an element in S
  - DELETE: Delete an element from S
- We can obviously sort using a priority queue:
  - Insert all elements using INSERT
  - Delete all elements in order using FINDMIN and DELETEMIN
- Priority queues have many other applications, e.g. in discrete event simulation, graph algorithms

2.1 Array or List implementations

- The first implementation that comes to mind is ordered array:

1	3	5	6	7	8	9	11	12	15	17

  - FINDMIN can be performed in $O(1)$ time
  - DELETEMIN and INSERT takes $O(n)$ time since we need to expand/compress the array after inserting or deleting element.
- If the array is unordered all operations take $O(n)$ time.

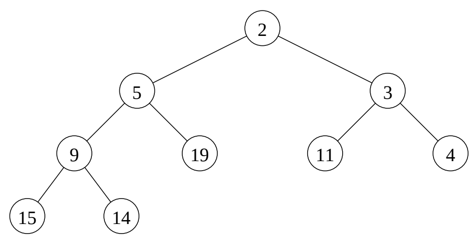
- We could use double linked sorted list instead of array to avoid the $O(n)$ expansion/compression cost
  - but INSERT can still take $O(n)$ time.

2.2 Heap implementation

- One way of implementing a priority queue is using a heap
- Heap definition:

  - Perfectly balanced binary tree
    - lowest level can be incomplete (but filled from left-to-right)
  - For all nodes v we have key(v)$\geq$key(parent(v))

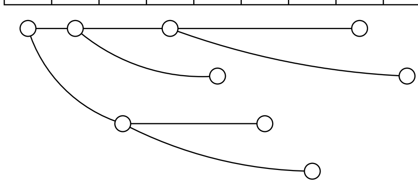
- Example:



- Heap can be implemented (stored) in two ways (at least)
  - Using pointers
  - In an array level-by-level, left-to-right

    Example:

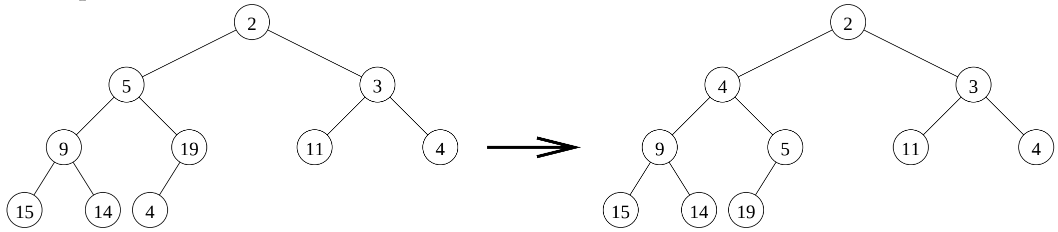
2	5	3	9	19	11	4	15	14



    - Note the nice property that the left and right children of node stored in entry i is in entry $2i$ and $2i+1$, respectively

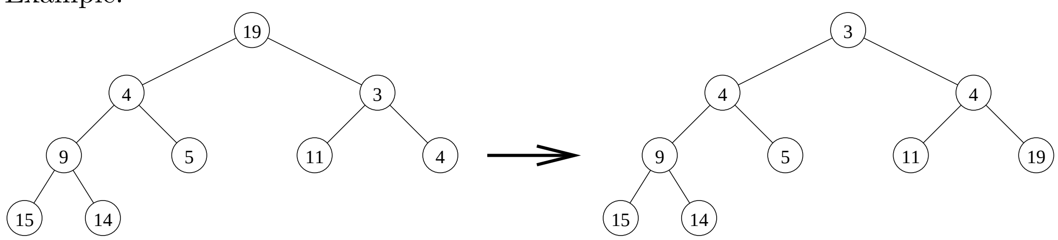
- Properties of heap:
  - Height $\Theta(\log n)$
  - Minimum of S is stored in root
- Operations:
  - INSERT
    - Insert element in new leaf in leftmost possible position on lowest level
    - Repeatedly swap element with element in parent node until heap order is reestablished (UP-HEAPIFY)

Example: Insertion of 4



- FindMin
  - Return root element
- DeleteMin
  - Delete element in root
  - Move element from rightmost leaf on lowest level to the root (and delete leaf)
  - Repeatedly swap element with element in child node with *minimal* element until heap order is reestablished (DOWN-HEAPIFY)

  Example:



- Running time: All operations traverse at most one root-leaf path $\Rightarrow O(\log n)$ time.
- Change and Delete can be handled similarly in $O(\log n)$ time
  - Assuming that we know the element to be changed/deleted.
- Sorting using heap (*Heap-Sort*) takes $\Theta(n \log n)$ time.
  - $n \cdot O(\log n)$ time to insert all elements (build the heap)
  - $n \cdot O(\log n)$ time to output sorted elements
- Sometimes we would like to build a heap faster than $O(n \log n)$
  - Insert elements in any order in perfectly balanced tree
  - DOWN-HEAPIFY all nodes level-by-level, bottom-up

  Correctness:

  - Induction on height of tree: When doing level i, all trees rooted at level $i-1$ are heaps.

  Analysis:

  - Define leaves to be on level 1 (root on level $\log n$)
  - n elements $\Rightarrow \leq \lceil \frac{n}{2} \rceil$ leaves $\Rightarrow \lceil \frac{n}{2^h} \rceil$ elements on level h
  - Cost of DOWN-HEAPIFY on a node on level h is h
  - Total cost: $\sum_{i=1}^{\log n} h \cdot \lceil \frac{n}{2^h} \rceil = \Theta(n) \cdot \sum_{i=1}^{\log n} \frac{h}{2^h}$
  - $\sum_{i=1}^{\log n} \frac{h}{2^h} = O(1)$ so cost is $\Theta(n)$
    - Assume $|x| < 1$ and differentiate $\sum_{h=0}^{\infty} x^h = \frac{1}{x-1}$
    - $\sum_{h=0}^{\infty} h x^{h-1} = \frac{1}{(x-1)^2} \Rightarrow \sum_{h=0}^{\infty} h x^h = \frac{x}{(x-1)^2} \Rightarrow \sum_{h=0}^{\infty} \frac{h}{2^h} = \frac{1/2}{(1/2-1)^2} = O(1)$