Algorithms Lab 11 (MST) csci 2200, Laura Toma. Bowdoin College

This week's topics

• MST: properties, the cut theorem; Kruskal's and Prim's algorithms

Review the algorithms discussed in class this week. This is a great time to understand all details and ask questions.

In lab: (COLLABORATION LEVEL 0: EVERYTHING ALLOWED)

- 1. **Prim's algorithm:** Following yesterday's discussion, we need to store the edges with one endpoint in T so that we can select the minimum-weight edge fast. We'll use a priority queue. The immediate idea is for the priority queue to store *edges*. Note however that for any vertex v not in T, even though it may have more than one edge connecting it to T, we'll only use the **smallest** one. So why store all? Just store the minimum edge connecting a vertex to T. This gives the idea for the priority queue to store *vertices*: each vertex v not in T yet has priority equal to the weight of the smallest edge that connectes it to a vertex u in T.
 - The priority queue will store all the vertices that are not in T yet.
 - A vertex v in the PQ has priority equal to the weight of the minimum edge that connects v with a vertex already in T. The other endpoint of this edge is stored in visit(v).
 - Essentially the priority queue stores all edges that *cross* the cut between the vertices in T and the vertices in V T (If a vertex not in T is connected by several edges to vertices in T, the pq will store only one of these edges, the smallest).
 - Initially T is empty and PQ contains all vertices in G with priority ∞ , except an arbitrary vertex v which has priority 0.
 - (a) Show how Prim's algorithm run on the example graph in the notes (or textbook).

- (b) What is the role of checking whether $v \in PQ$?
- (c) How many INSERT operations are performed by the algorithm?
- (d) How many DELETE-MIN operations are performed by the algorithm?
- (e) How many DECREASE-KEY operations are performed by the algorithm?
- (f) Assuming the priority is implemented as a heap, what is the complexity of the algorithm?
- 2. (CLRS 23.1-1) Consider an undircted, connected, weighted graph G. Is it true that the minimum-weight edge in G belongs to some MST of G? If yes, why?
- 3. (CLRS 24.2-4) Suppose that all edge weights in a graph are integers in the range from 1 to |V|. How can you take advantage of this in Kruskal's algorithm, and how fast can you make it run? What if the edge weights are integers from 1 to W for some constant W?
- 4. Is the path between a pair of nodes u and v in a minimum spanning tree (i.e. the path composed of edges in the MST) necessarily the minimum weight path between u and v if we consider all the edges in the graph G? If yes, prove it. If no, provide a counterexample.
- 5. Consider an undirected weighted graph which is formed by taking a binary tree and adding an edge from *exactly one* of the leaves to another node in the tree. We call such a graph a *loop-tree*. An example of a loop-tree could be the following:



Let n be the number of vertices in a loop-tree and assume that the graph is given in the normal edge-list representation without any extra information. In particular, the representation does not contain information about which vertex is the root.

- (a) How many edges are in a graph of n vertices?
- (b) How long time would it take Prim's or Kruskal's algorithms to find the minimal spanning tree of a loop-tree?
- (c) Describe and analyze a more efficient algorithm for finding the minimal spanning tree of a loop-tree. Remember to argue that the algorithm is correct.