

csci 210: Data Structures  
More Recursion

# Summary

- Topics: more recursion
  - Subset sum: finding if a subset of an array that sum up to a given target
  - Permute: finding all permutations of a given string
  - Subset: finding all subsets of a given string

# Thinking recursively

- Finding the recursive structure of the problem is the hard part
- Common patterns
  - divide in half, solve one half
  - divide in sub-problems, solve each sub-problem recursively, “merge”
  - solve one or several problems of size  $n-1$
  - process first element, recurse on remaining problem
- Recursion
  - functional: function computes and returns result.
    - Example: computing the sum of  $n$  numbers; isPalindrome; binary search.
  - procedural: no return result (function returns void). The task is accomplished during the recursive calls.
    - Example: Sierpinski fractals.
- Recursion
  - exhaustive
  - non-exhaustive: stops early

# Subset Sum

- Given an array of numbers and a target value, find whether there exists a subset of those numbers that sum up to the target value.

`boolean subsetSum (int[] a, int target)`

- Example:
- Recursive structure:
  - consider the next element in the array
  - try making a sum WITH this element
  - try making a sum WITHOUT this element
  - if neither is possible, return false

# Subset Sum

- So: consider the next element, it is either in the solution, or not. Try both ways. If both fail, return false.
- Need to keep track of the partial sum so far. When starting a recursive call, need to know the sum of the current subset. Also need to know the index of the next element to consider.

```
void recSubset(int[] a, int target, int i, int sumSoFar)
```

- The problem asked for a subsetSum function with the following signature:

```
boolean subsetSum (int[] a, int target)
```

- Need a wrapper:

```
boolean subsetSum (int[] a, int target) {  
    return recSubset(a, target, 0, 0);  
}
```

# Subset Sum

```
//i is the index of the next element to consider
//sumSoFar is the sum of elements included in the solution so far.
boolean recSubset(int[] a, int target, int i, int sumSoFar) {
    //basecases
    //we got it
    if (sumSoFar == target) return true;
    //we reached the end and sum is not equal to target
    if (i == a.length) return false;

    //recursive case: try next element both in and out of the sum
    boolean with = recSubset(a, target, i+1, sumSoFar + a[i]);
    boolean without = recSubset(a, target, i+1, sumSoFar);
    return (with || without);
}
```

# Subset Sum

- The tree of recursive calls for `recSubset([1, 2, 3, 4], target, 0, 0)`

# Subset Sum

- Variations
  - Alternative strategy: at each step, chose one of the remaining element to be part of the subset and recurse on the remaining part.
  - How could you change the function so that it prints the elements of the subset that sum to target?
    - store partial subsets in another array
    - or print element at the end of recursive call
  - How could you change the function to report not only if such a subset exists, but to count all such subsets?



# Permutations

- Write a function to print all permutations of a given string.
- Example: permute “abc” should print: abc, acb, bca, bac, cab, cba.

```
void printPerm(String s)
```

- Recursive structure:
  - Chose a letter from the input, and make this the first letter of the output
  - Recursively permute remaining input
  - chose a, permute “bc”: should generate “a” + all permutations of “bc”
  - chose all letters in turn to be first letters
  - chose b, permute “ac”: should generate “b” + all permutations of “ac”
  - chose c, permute “ab”: should generate “c” + all permutations of “ab”
  
- What is the base case?
- Can you make sure that each permutation is generated precisely once?

# Permutations

- So: pick a letter, add it to the solution, recurse on remaining
- When starting a recursive call, we know the list of letters chosen so far; that is, we know the first part of the permutation generated so far.
- Need to keep track of it.

```
//print soFar + all permutations of remaining  
void recPermute(String soFar, String remaining)
```

- The problem asked for a printPermute with a different signature: we need a wrapper

```
//print all permutations of s  
void printPerm (String s) {  
    recPermute("", s);  
}
```

- Why use wrappers? the user does not need to know the internals of the implementation. In this case, that it is recursive.

# Permutations

```
void recPermute(String soFar, String remaining) {  
  
    //base case  
    if (remaining.length() == 0)  
        System.out.println(soFar);  
    else {  
        for (int i=0; i< remaining.length(); i++) {  
            String nextSoFar = soFar + remaining[i];  
            String nextRemaining = remaining.substring(0,i) + remaining.substring(i+1);  
            recPermute(nextSoFar, nextRemaining)  
        }  
    }  
}
```

# Permutations

- The tree of recursive calls for `recPermute("", "abc")`

# Subsets

- Enumerate all subsets of a given string
- Example: subsets of “abc” are a, b, c, ab, ac, bc, abc
  - Order does not matter: “ab” is the same as “ba”
- Recursive structure
  - chose one element from input
  - can either include it in current subset or not
  - recursively form subsets including it
  - recursively form subsets excluding it
  - make sure to generate each set once
  - base case?

# Subsets

```
void recSubsets(String soFar, String remaining) {  
    if (remaining.length()==0)  
        System.out.println(soFar);  
    else {  
        //add to subset, remove from rest, recurse  
        recSubsets(soFar+remaining[0], remaining.substring(1));  
        //don't add to subset, remove from rest, recurse  
        recSubsets(soFar, remaining.substring(1));  
    }  
}
```

```
void subsets(String s) {  
    recSubsets("", s);  
}
```

# Subsets

- The tree of recursive calls for `recSubsets("", "abcd")`