

csci 210: Data Structures

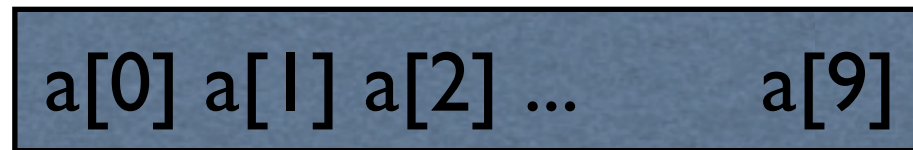
Linked lists

Summary

- Today
 - linked lists
 - single-linked lists
 - double-linked lists
 - circular lists
- **READING:**
 - GT textbook chapter 3.2. 3.3. 3.4

Arrays vs. Linked Lists

- We've seen arrays:
 - `int[] a = new int[10];`
 - a is a chunk of memory of size $10 \times \text{sizeof}(\text{int})$
 - a has a fixed size



- A linked list is fundamentally different way of storing collections
 - each element stores a reference to the element after it



Arrays vs. Lists

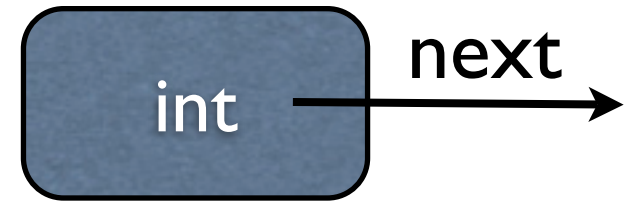
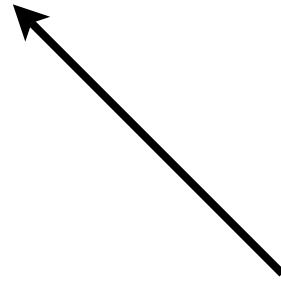
- **Arrays**
 - have a pre-determined fixed size
 - easy access to any element $a[i]$ in constant time
 - no space overhead
 - $\text{Size} = n \times \text{sizeof}(\text{element})$

- **Linked lists**
 - no fixed size; grow one element at a time
 - space overhead
 - each element must store an additional reference
 - $\text{Size} = n \times \text{sizeof}(\text{element}) + n \times \text{sizeof}(\text{reference})$
 - no easy access to i -th element wrt the head of the list
 - need to hop through all previous elements

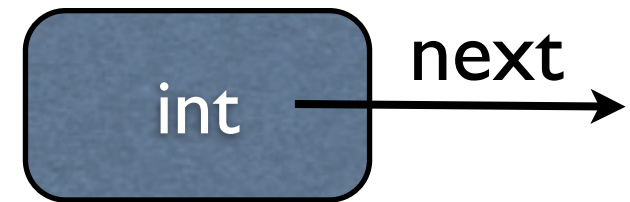
The Node class

```
/** Node of a singly linked list of integers */  
public class Node {  
  
    private int element; // we assume elements are ints  
    private Node next;
```

self-referential definition



The Node class



```
/** Node of a singly linked list of integers */
public class Node {
    private int element;    // we assume elements are ints
    private Node next;

    /** Creates a node with the given element and next node. */
    public Node(int s, Node n) {
        element = s;
        next = n;
    }

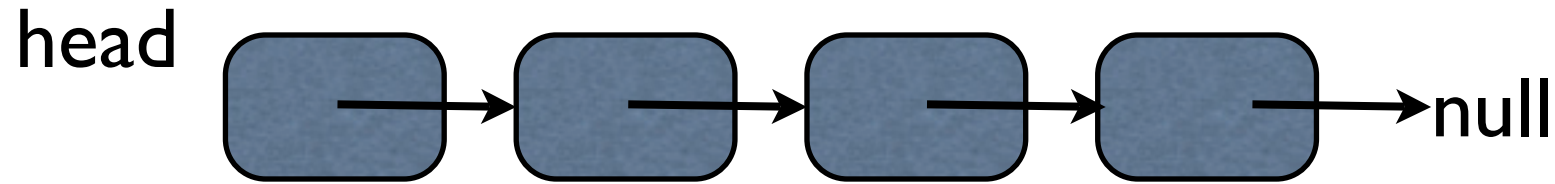
    /** Returns the element of this node. */
    public int getElement() { return element; }

    /** Returns the next node of this node. */
    public Node getNext() { return next; }

    // Modifier methods:
    /** Sets the element of this node. */
    public void setElement(int newElem) { element = newElem; }

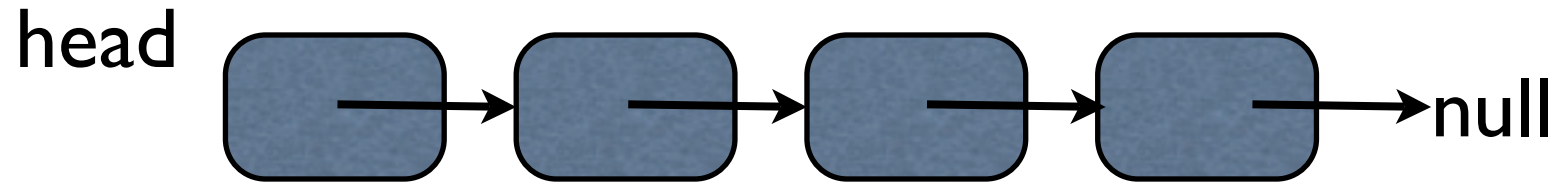
    /** Sets the next node of this node. */
    public void setNext(Node newNext) { next = newNext; }
}
```

A Single-Linked-List class



```
/** Singly linked list */  
public class SLinkedList {  
    protected Node head;        // head node of the list  
    protected long size;        // number of nodes in the list  
  
    /** Default constructor that creates an empty list */  
    public SLinkedList() {  
        head = null;  
        size = 0;  
    }  
}
```

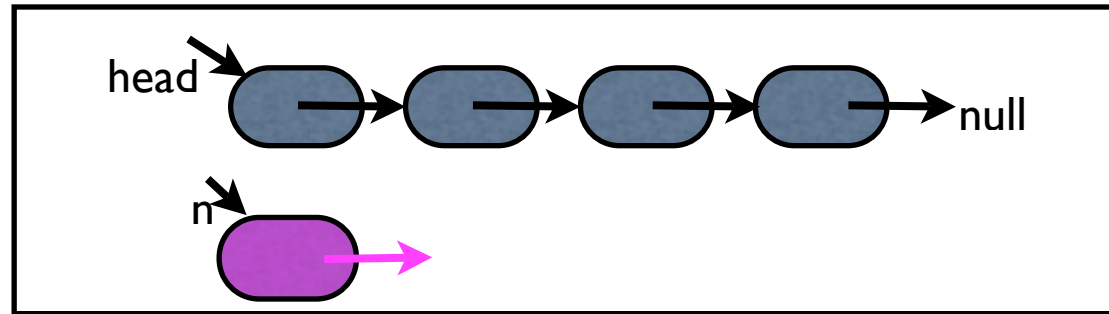
A Single-Linked-List class



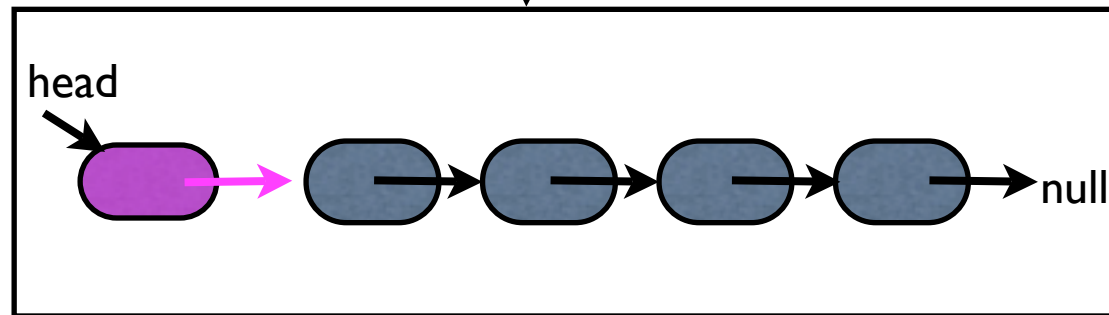
```
/** Singly linked list */  
public class SLinkedList {  
    protected Node head;        // head node of the list  
    protected long size;        // number of nodes in the list  
  
    /** Default constructor that creates an empty list */  
    public SLinkedList() {  
        head = null;  
        size = 0;  
    }  
  
}
```

- we'll discuss the following methods
 - addFirst(Node n)
 - addAfter(Node n)
 - Node get(int i)
 - Node removeFirst()
 - addLast(Node n)
 - removeLast(Node n)

Inserting at head



```
void addFirst(Node n)
```

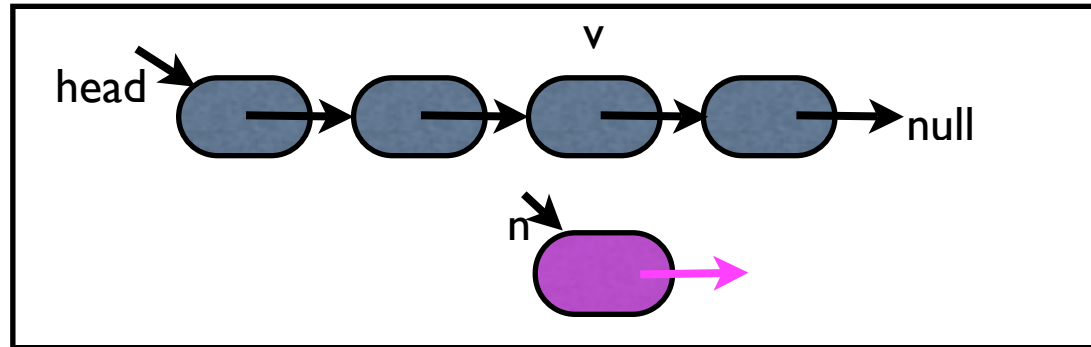


```
void addFirst(Node n) {  
    n.setNext(head);  
    head = n;  
    size++;  
}
```

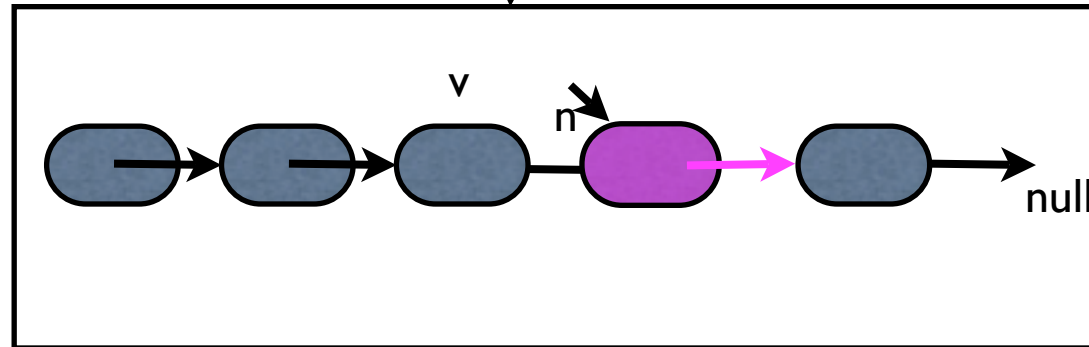
- Notes

- Special cases: works when head is null, i.e. list is empty
- Efficiency: $O(1)$ time (i.e. constant time)

Inserting in the middle



```
void insertAfter(Node v, Node n)
```



```
//insert node n after node v
```

```
void insertAfter(Node v, Node n)
```

```
    n.setNext(v.getNext());
```

```
    v.setNext(n);
```

```
    size++;
```

```
}
```

- Notes:

- Efficiency: $O(1)$ (constant time)
- Special cases
 - does not work if v or n are null
 - null pointer exception

Get the i-th element

```
//return the i-th node  
Node get(int i)
```

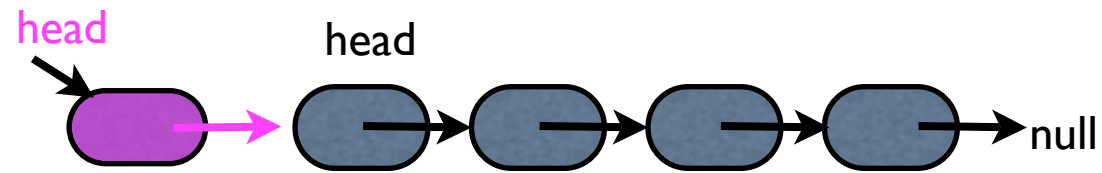
Get the i-th element

```
//return the i-th node
Node get(int i) {
    if (i >= size) print error message and return null
    Node ptr = head;
    for (int k=0; k<i; k++)
        ptr = ptr.getNext();
    return ptr;
}
```

- Notes

- Special cases
 - does it work when list is empty?
- Efficiency: takes $O(i)$ time
 - constant time per element traversed
 - unlike arrays, accessing i-th element is not constant time

Remove at head



```
Node removeFirst() {  
    Node n = head;  
    head = head.getNext();  
    n.setNext(null);  
    return n;  
}
```

- Notes:
 - Special cases
 - does it work when list is empty?
 - Nope.
 - How to fix it?
 - Efficiency: $O(1)$

Insert at tail

```
void addLast(Node n) {  
    insertAfter (get(size), n);  
}
```

- Notes

- Special cases
 - does it work when list is empty?
 - Nope (first node in insertAfter is null).
 - How to fix it?
- Efficiency: takes $O(\text{size})$ time

Delete at tail

- Remove at end: similar
 - need to get to the last element from the head
 - $O(\text{size})$ time

Linked lists

- Single-linked lists support insertions and deletions **at head** in $O(1)$ time.
 - insertions and deletion at the tail can be supported in $O(\text{size})$ time.
 - addFirst: $O(1)$ time (constant time)
 - removeFirst: $O(1)$ time (constant time)
 - addLast: $O(\text{size})$ time
 - removeLast: $O(\text{size})$ time
- Why? because we keep track of the head.
 - To access the tail in constant time, need to keep track of tail as well.

Linked-list with tail

```
/** Singly linked list */  
public class SLinkedList {  
  
    private Node head, tail; // head and tail nodes of the list  
    private long size;      // number of nodes in the list
```

```
void SLinkedList() {  
    head = tail = null;  
    size = 0;  
}  
  
void addFirst(Node n) {...}  
  
Node removeFirst() {...}  
  
....  
}
```

← all methods must update tail

Insert at tail

- Special cases: list is empty

```
void addLast(Node n) {  
  
    //if list is empty the new element is head and tail  
    if (tail == null) {  
        n.setNext(null);  
        head = tail = n;  
    } else {  
        //the list is not empty: link tail to n and n becomes the new tail  
        tail.setNext(n);  
        n.setNext(null);  
        tail = n;  
    }  
    //increment size  
    size++  
}
```

- Efficiency: $O(1)$

Remove at tail

- What we want: delete the last element and set the new tail
- Is that possible?

Remove at tail

- What we want: delete the last element and set the new tail
- Is that possible?

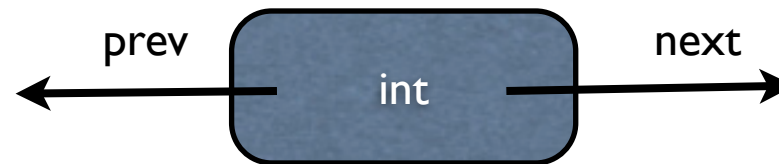
- Remove at tail
 - set the tail to the node BEFORE the tail
 - need the node before the tail: $O(\text{size})$

- To remove an element from a list you need the node BEFORE it as well

```
remove(Node n) {  
    //link n.before to n.next  
}
```

- To remove a node efficiently need to keep track of previous node

Doubly-linked lists



```
/** Node of a doubly linked list of integers */
public class DNode {
    protected int element;    //element stored by a node
    protected DNode next, prev;    // Pointers to next and previous nodes

    /** Constructor that creates a node with given fields */
    public DNode(int e, DNode p, DNode n) {
        element = e;
        prev = p;
        next = n;
    }
    /** Returns the element of this node */
    public int getElement() { return element; }
    /** Returns the previous node of this node */
    public DNode getPrev() { return prev; }
    /** Returns the next node of this node */
    public DNode getNext() { return next; }
    /** Sets the element of this node */
    public void setElement(int newElem) { element = newElem; }
    /** Sets the previous node of this node */
    public void setPrev(DNode newPrev) { prev = newPrev; }
    /** Sets the next node of this node */
    public void setNext(DNode newNext) { next = newNext; }
}
```

Doubly-linked lists

```
/** Doubly linked list with nodes of type DNode */
public class DList {

    protected int size;           // number of elements
    protected DNode head, tail;

    void addFirst(Node n);
    void addLast(Node n);
    Node deleteFirst();
    Node deleteLast();
    void delete(Node n);
}
```

- Operations on doubly linked lists
 - addFirst(): $O(1)$ time
 - addLast(): $O(1)$ time
 - deleteFirst(): $O(1)$ time
 - deleteLast(): $O(1)$ time
 - delete(): $O(1)$ time
 - get(i): $O(i)$ time

Insert at head

```
void addFirst(Node n) {  
  
    n.setNext(head);  
    n.setprev(null);  
    head.setPrev(n);  
    head = n;  
    size++;  
}
```

Does this work?

Insert at head

```
void addFirst(Node n) {
    n.setNext(head);
    n.setprev(null);
    head.setPrev(n);
    head = n;
    size++;
}
```

- Special cases?

- empty list: head is null; need to set tail too

```
void addFirst(Node n) {
    if (head==null) {
        //this is the first element: set both head and tail to it
        head = tail = n;
        n.setPrev(null); n.setNext(null);
    }
    else {
        n.setNext(head); n.setprev(null);
        head.setPrev(n);
        head = n;
    }
    size++;
}
```

- Efficiency: $O(1)$

Insert at tail

```
void addLast(Node n) {  
    tail.setNext(n);  
    n.setprev(tail);  
    n.setNect(null);  
    tail = n;  
    size++;  
}
```

Does this work ?

Insert at tail

```
void addLast(Node n) {
    tail.setNext(n);
    n.setprev(tail);
    n.setNect(null);
    tail = n;
    size++;
}
```

- Special cases?
 - empty list: tail is null; need to set head too

```
void addLast(Node n) {
    if (tail == null) {
        head = tail = n; n.setPrev(null); n.setNext(null);
    }
    else {
        tail.setNext(n); n.setprev(tail); n.setNect(null);
        tail = n;
    }
    size++;
}
```

- Efficiency: $O(1)$

Doubly-linked lists

- Class work: Sketch the following methods for doubly-linked lists, and analyze their efficiency.
 - Node removeFirst()
 - Node removeLast()
 - void remove(Node n)
 - Node search(int k)

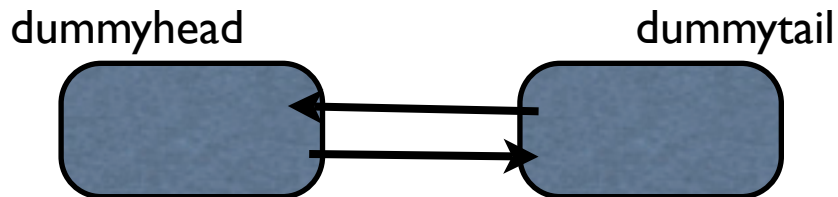
Sentinels

- Sentinels for singly-linked list: keep a dummy head
 - an empty list is one node: the dummy head
- Sentinels for doubly-linked lists
 - dummy head and dummy tail
- Why? elegant. Unifies special cases when head or tail are null

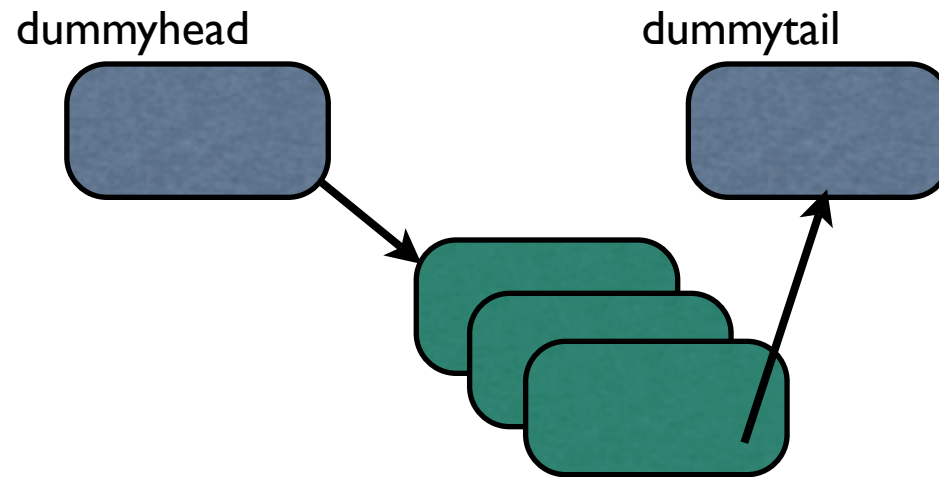
DLLists with Sentinels

```
public class DList {  
  
    protected int size;           // number of elements  
    protected DNode header, trailer; // sentinels  
  
    /** Constructor that creates an empty list */  
    public DList() {  
        size = 0;  
        header = new DNode(null, null, null); // create header  
        trailer = new DNode(null, header, null); // create trailer  
        // make header and trailer point to each other  
        header.setNext(trailer);  
  
    }  
}
```

- the empty list:
 - size = 0

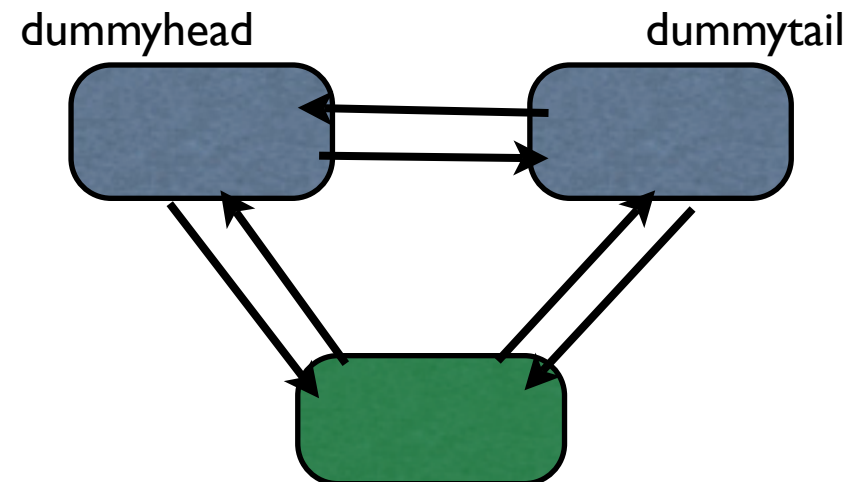


DLLists with sentinels



```
insertFirst(Node n) {  
    n.setNext(dummyHead.getNext());  
    dummyHead.getNext().setPrev(n);  
    dummyHead.setNext(n);  
    n.setPrev(dummyhead);  
    size++;  
}
```

- Special cases: none
 - works for empty list



Extensions

- Circular lists
 - make last node point to the first (instead of null)

- `class CircularList {`

- `SNode head;`
 - `int size;`

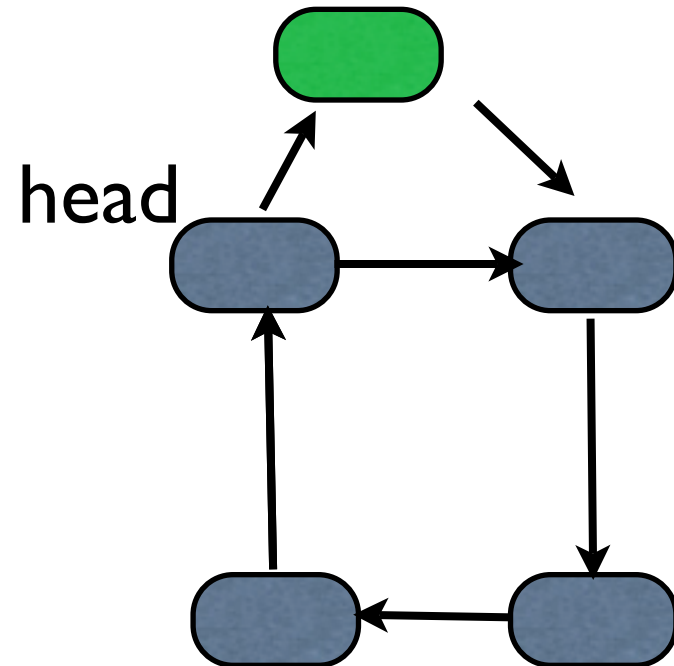
- `}`

- let's say we want to insert at head

```
insertAtHead(Node n) {  
    n.setNext(head.getNext());  
    head.setNext(n);  
}
```

- if head is null?

```
if (head == null) {  
    n.setNext(n);  
    head = n;  
}
```



Linked-lists in Java

- search for class Java Linked List
- has all expected methods and features