

csci 210: Data Structures

Linked lists

Summary

- Today
 - linked lists
 - single-linked lists
 - double-linked lists
 - circular lists
- READING:
 - LC chapter 4.1, 4.2, 4.3

Arrays vs. Linked Lists

- We've seen arrays:
 - `int[] a = new int[10];`
 - a is a chunk of memory of size $10 \times \text{sizeof}(\text{int})$
 - a has a fixed size

a[0] a[1] a[2] ... a[9]

- A linked list is fundamentally different way of storing collections
 - each element stores a reference to the element after it



Arrays vs. Lists

- Arrays
 - have a pre-determined fixed size
 - easy access to any element $a[i]$ in constant time
 - no space overhead
 - $\text{Size} = n \times \text{sizeof}(\text{element})$
- Linked lists
 - no fixed size; grow one element at a time
 - space overhead
 - each element must store an additional reference
 - $\text{Size} = n \times \text{sizeof}(\text{element}) + n \times \text{sizeof}(\text{reference})$
 - no easy access to i -th element wrt the head of the list
 - need to hop through all previous elements

LinkedLists in Java

- Search for class Java LinkedList
- Has all expected methods and features
 - add(int index, Object element)
 - add(Object o)
 - addAll(Collection c)
 - addAll(int index, Collection c)
 - addFirst(Object o)
 - addLast(Object o)
 - contains(Object o)
 - get(int index)
 - getFirst()
 - getLast()
 - indexOf(Object o)
 - lastIndexOf(Object o)
 - remove(int index)
 - remove(Object o)
 - removeFirst()
 - removeLast()
 - set(int index, Object element)
 - size()

Implementing a linked list

- We want to implement a linked list class, much like Java's LinkedList
- For simplicity, we can think of a linked list of integers

The Node class



We want to define the node in a list linked of integers.

```
/** Node of a singly linked list of integers */
public class Node {
    ...
}
```

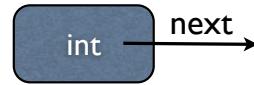
The Node class



We want to define the node in a list linked of integers.

```
/** Node of a singly linked list of integers */
public class Node {
    private int element; //we assume elements are ints
    private Node next;
    ...
}
self-referential definition
```

The Node class



```
/** Node of a singly linked list of integers */
public class Node {
    private int element; // we assume elements are ints
    private Node next;

    /** Creates a node with the given element and next node. */
    public Node(int s, Node n) {
        element = s;
        next = n;
    }

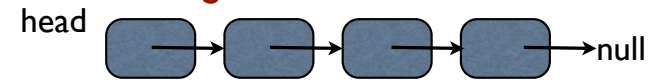
    /** Returns the element of this node. */
    public int getElement() { return element; }

    /** Returns the next node of this node. */
    public Node getNext() { return next; }

    // Modifier methods:
    /** Sets the element of this node. */
    public void setElement(int newElem) { element = newElem; }

    /** Sets the next node of this node. */
    public void setNext(Node newNext) { next = newNext; }
}
```

A Single-Linked-List class

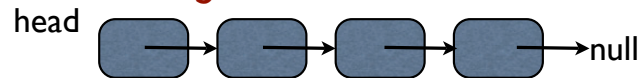


```
/** Singly linked list */
public class SLinkedList {
    protected Node head; // head node of the list
    protected long size; // number of nodes in the list

    /** Default constructor that creates an empty list */
    public SLinkedList() {
        head = null;
        size = 0;
    }

    ...
}
```

A Single-Linked-List class



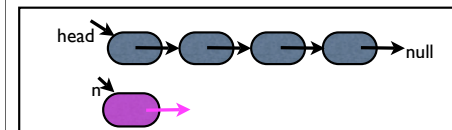
```
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    /** Default constructor that creates an empty list */
    public SLinkedList() {
        head = null;
        size = 0;
    }

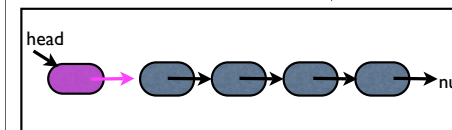
    ...
}
```

- We'll discuss the following methods
 - addFirst(Node n)
 - addAfter(Node n)
 - Node get(int i)
 - Node removeFirst()
 - addLast(Node n)
 - removeLast(Node n)

Inserting at head



```
void addFirst(Node n)
```

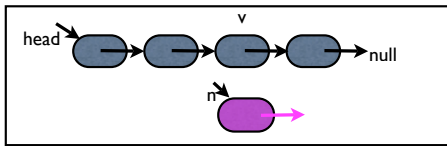


```
void addFirst(Node n) {
    n.setNext(head);
    head = n;
    size++;
}
```

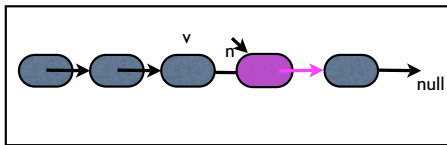
• Notes

- Special cases: works when head is null, i.e. list is empty
- Efficiency: $O(1)$ time

Inserting in the middle



```
void insertAfter(Node v, Node n)
```



```
//insert node n after node v  
void insertAfter(Node v, Node n)  
{  
    n.setNext(v.getNext());  
    v.setNext(n);  
    size++;  
}
```

• Notes:

- Efficiency: $O(1)$
- Special cases
 - does not work if v or n are null
 - null pointer exception

Get the i-th element

```
//return the i-th node  
Node get(int i) {  
    ...  
  
}
```

Get the i-th element

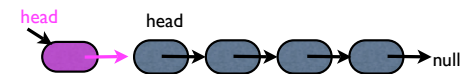
```
//return the i-th node
```

```
Node get(int i) {  
    if (i >= size) print error message and return null  
    Node ptr = head;  
    for (int k=0; k<i; k++)  
        ptr = ptr.getNext();  
    return ptr;  
}
```

• Notes

- Special cases
 - does it work when list is empty?
- Efficiency: takes $O(i)$ time
 - constant time per element traversed
 - unlike arrays, accessing i-th element is not constant time

Remove at head



```
Node removeFirst() {  
    Node n = head;  
    head = head.getNext();  
    n.setNext(null);  
    return n;  
}
```

Notes:

- Special cases
 - does it work when list is empty?
 - Nope.
 - How to fix it?
- Efficiency: $O(1)$

Insert at tail

```
void addLast(Node n) {  
    insertAfter (get(size), n);  
}
```

- Notes
- Special cases
 - does it work when list is empty?
 - Nope (first node in insertAfter is null).
 - How to fix it?
- Efficiency: takes $O(\text{size})$ time


Delete at tail

- Remove at end: similar
 - need to get to the last element from the head
 - $O(\text{size})$ time

Linked lists

- Single-linked lists support insertions and deletions **at head** in $\Theta(1)$ time.
- Insertions and deletion at the tail can be supported in $O(\text{size})$ time.
 - addFirst: $O(1)$ time
 - removeFirst: $O(1)$ time
 - addLast: $O(\text{size})$ time
 - removeLast: $O(\text{size})$ time
- Why? because we keep track of the head.
 - To access the tail in constant time, need to keep track of tail as well.

Linked-list with tail

```
/** Singly linked list .*/  
public class SLinkedList {  
  
    private Node head, tail; // head and tail nodes of the list  
    private long size;      // number of nodes in the list  
  
    void SLinkedList() {  
        head = tail = null;  
        size = 0;  all methods must update tail  
    }  
  
    void addFirst(Node n) {...}  
  
    Node removeFirst() {...}  
  
    ....  
}
```

Insert at tail

```
void addLast(Node n) {  
    //if list is empty the new element is head and tail  
    if (tail == null) {  
        n.setNext(null);  
        head = tail = n;  
    } else {  
        //the list is not empty: link tail to n and n becomes the new  
        //tail  
        tail.setNext(n);  
        n.setNext(null);  
        tail = n;  
    }  
    //increment size  
    size++;  
}
```

- Special cases: list is empty
- Efficiency: Theta(1)

Remove at tail

- What we want: delete the last element and set the new tail
- Is that possible?

Remove at tail

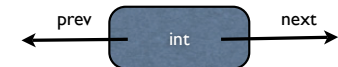
- What we want: delete the last element and set the new tail
- Is that possible?

- Remove at tail
 - set the tail to the node BEFORE the tail
 - need the node before the tail: O(size)
- To remove an element from a list you need the node BEFORE it as well

```
remove(Node n) {  
    //link n.before to n.next  
}
```

- To remove a node efficiently need to keep track of `previous node`

Doubly-linked lists



```
/** Node of a doubly linked list of integers */  
public class DNode {  
    protected int element; //element stored by a node  
    protected DNode next, prev; // Pointers to next and previous nodes  
  
    /** Constructor that creates a node with given fields */  
    public DNode(int e, DNode p, DNode n) {  
        element = e;  
        prev = p;  
        next = n;  
    }  
    /** Returns the element of this node */  
    public int getElement() { return element; }  
    /** Returns the previous node of this node */  
    public DNode getPrev() { return prev; }  
    /** Returns the next node of this node */  
    public DNode getNext() { return next; }  
    /** Sets the element of this node */  
    public void setElement(int newElem) { element = newElem; }  
    /** Sets the previous node of this node */  
    public void setPrev(DNode newPrev) { prev = newPrev; }  
    /** Sets the next node of this node */  
    public void setNext(DNode newNext) { next = newNext; }  
}
```

Doubly-linked lists

```
/** Doubly linked list with nodes of type DNode */
public class DList {

    protected int size; // number of elements
    protected DNode head, tail;

    void addFirst(Node n);
    void addLast(Node n);
    Node deleteFirst();
    Node deleteLast();
    void delete(Node n);
}
```

- Operations on doubly linked lists
 - addFirst(): O(1) time
 - addLast(): O(1) time
 - deleteFirst(): O(1) time
 - deleteLast(): O(1) time
 - delete(): O(1) time
 - get(i): O(i) time

Insert at head

```
void addFirst(Node n) {
    n.setNext(head);
    n.setprev(null);
    head.setPrev(n);
    head = n;
    size++;
}
```

Does this work?

Insert at head

```
void addFirst(Node n) {
    n.setNext(head);
    n.setprev(null);
    head.setPrev(n);
    head = n;
    size++;
}

void addFirst(Node n) {
    if (head==null) {
        /* this is the first
        element: set both head
        and tail to it */
        head = tail = n;
        n.setPrev(null);
        n.setNext(null);
    }
    else {
        n.setNext(head);
        n.setprev(null);
        head.setPrev(n);
        head = n;
    }
    size++;
}
```

- Special cases?
 - empty list: head is null; need to set tail too
- Efficiency ?
 - O(1)

Insert at tail

```
void addLast(Node n) {
    tail.setNext(n);
    n.setprev(tail);
    n.setNect(null);
    tail = n;
    size++;
}
```

Does this work ?

Insert at tail

```
void addLast(Node n) {
    tail.setNext(n);
    n.setprev(tail);
    n.setNect(null);
    tail = n;
    size++;
}

void addLast(Node n) {
    if (tail == null) {
        head = tail = n;
        n.setPrev(null);
        n.setNext(null);
    }
    else {
        tail.setNext(n);
        n.setprev(tail);
        n.setNect(null);
        tail = n;
    }
    size++;
}
```

- Special cases?
 - empty list: tail is null; need to set head too
- Efficiency: $O(1)$

Doubly-linked lists

- Class work: Sketch the following methods for doubly-linked lists, and analyze their efficiency.

- Node removeFirst()
- Node removeLast()
- void remove(Node n)
- Node search(int k)

Sentinels

- Sentinels for singly-linked list: keep a dummy head
 - an empty list is one node: the dummy head
- Sentinels for doubly-linked lists
 - dummy head and dummy tail
- Why? elegant. Unifies special cases when head or tail are null

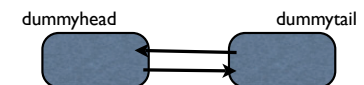
DLLists with Sentinels

```
public class DList {
    protected int size; // number of elements
    protected DNode header, trailer; // sentinels

    /** Constructor that creates an empty list */
    public DList() {
        size = 0;
        header = new DNode(null, null, null); // create header
        trailer = new DNode(null, header, null); // create trailer
        // make header and trailer point to each other
        header.setNext(trailer);
    }
}
```

- the empty list:

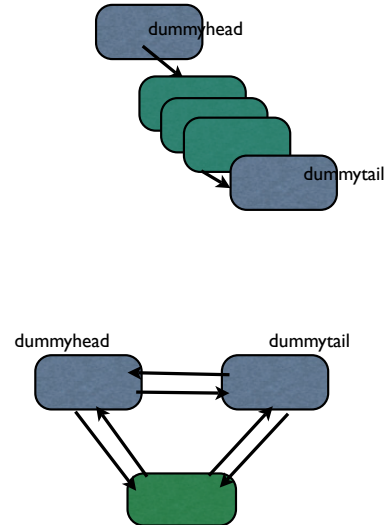
- size = 0



DLLists with sentinels

```
insertFirst(Node n) {  
  n.setNext(dummyHead.getNext());  
  dummyHead.getNext().setPrev(n);  
  dummyHead.setNext(n);  
  n.setPrev(dummyhead);  
  size++;  
}
```

- Special cases: none
 - works for empty list



Extensions

- Circular lists : make last node point to the first (instead of null)

```
class CircularList {  
  SNode head;  
  int size;  
}  
  
insertAtHead(Node n) {  
  n.setNext(head.getNext());  
  head.setNext(n);  
}  
  
if (head == null) {  
  n.setNext(n);  
  head = n;  
}
```

- Let's say we want to insert at head

