

csci 210: Data Structures

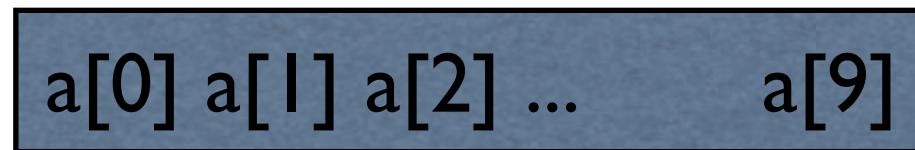
Linked lists

Summary

- Today
 - linked lists
 - single-linked lists
 - double-linked lists
 - circular lists
- READING:
 - LC chapter 4.1, 4.2, 4.3

Arrays vs. Linked Lists

- We've seen arrays:
 - `int[] a = new int[10];`
 - `a` is a chunk of memory of size $10 \times \text{sizeof}(\text{int})$
 - `a` has a fixed size



- A linked list is fundamentally different way of storing collections
 - each element stores a reference to the element after it



Arrays vs. Lists

- **Arrays**
 - have a pre-determined fixed size
 - easy access to any element $a[i]$ in constant time
 - no space overhead
 - $\text{Size} = n \times \text{sizeof(element)}$
- **Linked lists**
 - no fixed size; grow one element at a time
 - space overhead
 - each element must store an additional reference
 - $\text{Size} = n \times \text{sizeof(element)} + n \times \text{sizeof(reference)}$
 - no easy access to i -th element wrt the head of the list
 - need to hop through all previous elements

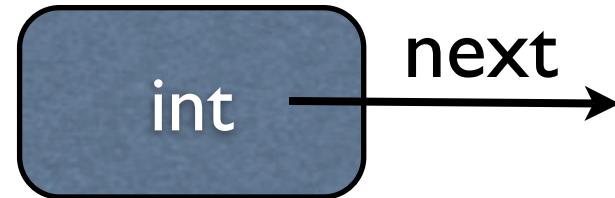
Linked-lists in Java

- Search for class Java LinkedList
- Has all expected methods and features
 - add(int index, Object element)
 - add(Object o)
 - addAll(Collection c)
 - addAll(int index, Collection c)
 - addFirst(Object o)
 - addLast(Object o)
 - contains(Object o)
 - get(int index)
 - getFirst()
 - getLast()
 - indexOf(Object o)
 - lastIndexOf(Object o)
 - remove(int index)
 - remove(Object o)
 - removeFirst()
 - removeLast()
 - set(int index, Object element)
 - size()

Implementing a linked list

- We want to implement a linked list class, much like Java's `LinkedList`
- For simplicity, we can think of a linked list of integers

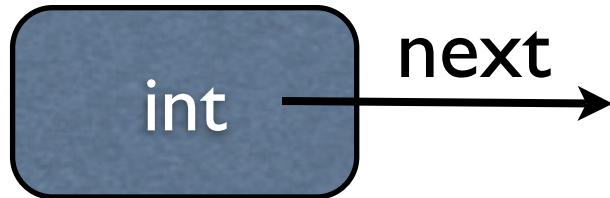
The Node class



We want to define the node in a list linked of integers.

```
/** Node of a singly linked list of integers */
public class Node {  
    ...  
}
```

The Node class



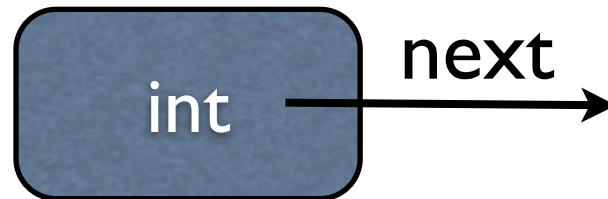
We want to define the node in a list linked of integers.

```
/** Node of a singly linked list of integers */
public class Node {  
  
    private int element; //we assume elements are ints  
    private Node next;  
  
    ...  
  
}
```

self-referential definition

An arrow originates from the "next" field in the code and points to the "next" label in the diagram above, highlighting the self-referential nature of the definition.

The Node class



```
/** Node of a singly linked list of integers */
public class Node {
    private int element; // we assume elements are ints
    private Node next;

    /** Creates a node with the given element and next node. */
    public Node(int s, Node n) {
        element = s;
        next = n;
    }

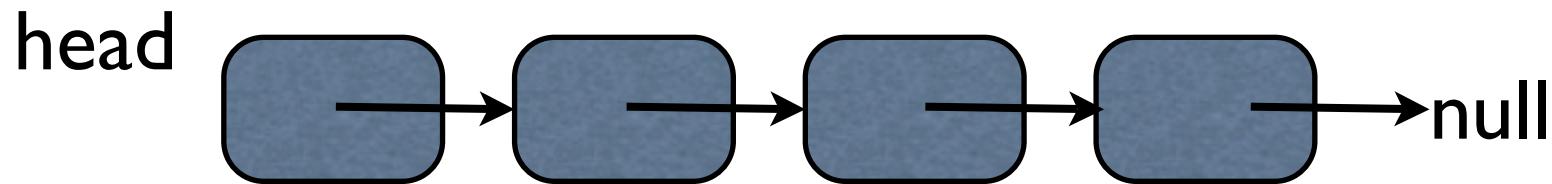
    /** Returns the element of this node. */
    public int getElement() { return element; }

    /** Returns the next node of this node. */
    public Node getNext() { return next; }

    // Modifier methods:
    /** Sets the element of this node. */
    public void setElement(int newElem) { element = newElem; }

    /** Sets the next node of this node. */
    public void setNext(Node newNext) { next = newNext; }
}
```

A Single-Linked-List class

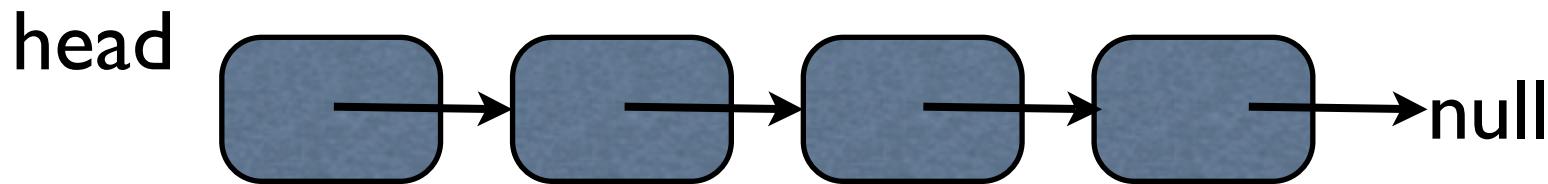


```
/** Singly linked list */
public class SLinkedList {
    protected Node head;          // head node of the list
    protected long size;          // number of nodes in the list

    /** Default constructor that creates an empty list */
    public SLinkedList() {
        head = null;
        size = 0;
    }

    ...
}
```

A Single-Linked-List class

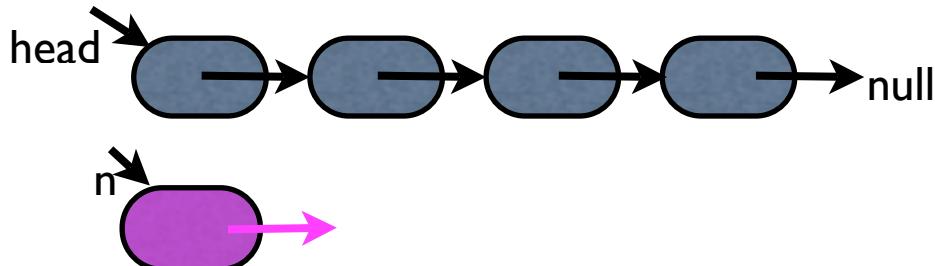


```
/** Singly linked list.*/
public class SLinkedList {
    protected Node head;      // head node of the list
    protected long size;       // number of nodes in the list

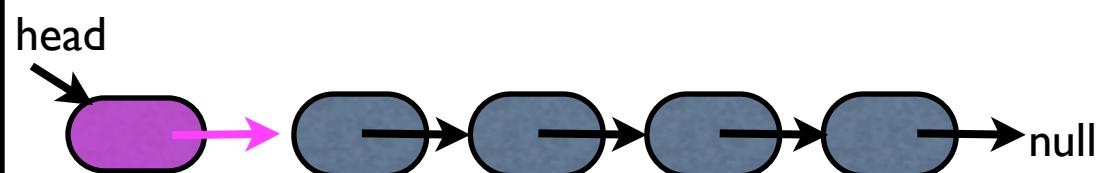
    /** Default constructor that creates an empty list */
    public SLinkedList() {
        head = null;
        size = 0;
    }
    ...
}
```

- We'll discuss the following methods
 - `addFirst(Node n)`
 - `addAfter(Node n)`
 - `Node get(int i)`
 - `Node removeFirst()`
 - `addLast(Node n)`
 - `removeLast(Node n)`

Inserting at head



```
void addFirst(Node n)
```

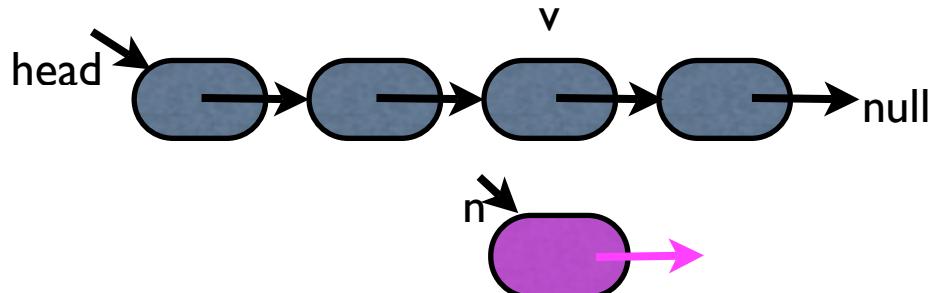


```
void addFirst(Node n) {  
    n.setNext(head);  
    head = n;  
    size++;  
}
```

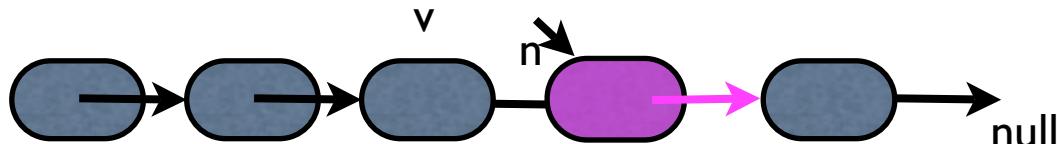
- Notes

- Special cases: works when head is null, i.e. list is empty
- Efficiency: O(1) time

Inserting in the middle



```
void insertAfter(Node v, Node n)
```



```
//insert node n after node v  
void insertAfter(Node v, Node n)  
    n.setNext(v.getNext());  
    v.setNext(n);  
    size++;  
}
```

- Notes:

- Efficiency: O(1)
- Special cases
 - does not work if v or n are null
 - null pointer exception

Get the i-th element

```
//return the i-th node  
  
Node get(int i) {  
  
    ...  
  
}
```

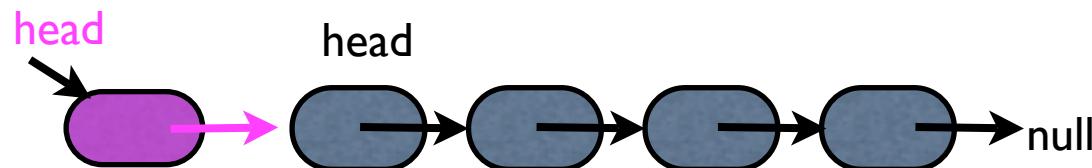
Get the i-th element

```
//return the i-th node

Node get(int i) {
    if (i >= size) print error message and return null
    Node ptr = head;
    for (int k=0; k<i; k++)
        ptr = ptr.getNext();
    return ptr;
}
```

- Notes
 - Special cases
 - does it work when list is empty?
 - Efficiency: takes $O(i)$ time
 - constant time per element traversed
 - unlike arrays, accessing i-th element is not constant time

Remove at head



```
Node removeFirst() {  
    Node n = head;  
    head = head.getNext();  
    n.setNext(null);  
    return n;
```

Notes:

- Special cases
 - does it work when list is empty?
 - Nope.
 - How to fix it?
- Efficiency: O(1)

Insert at tail

```
void addLast(Node n) {  
    insertAfter (get(size), n);  
}
```

- Notes
- Special cases
 - does it work when list is empty?
 - Nope (first node in insertAfter is null).
 - How to fix it?
 - Efficiency: takes O(size) time

Delete at tail

- Remove at end: similar
 - need to get to the last element from the head
 - $O(\text{size})$ time

Linked lists

- Single-linked lists support insertions and deletions **at head** in Theta(1) time.
- Insertions and deletion at the tail can be supported in O(size) time.
 - addFirst: O(1) time
 - removeFirst: O(1) time
 - addLast: O(size) time
 - removeLast: O(size) time
- Why? because we keep track of the head.
 - To access the tail in constant time, need to keep track of tail as well.

Linked-list with tail

```
/** Singly linked list */
public class SLinkedList {

    private Node head, tail; // head and tail nodes of the list
    private long size;        // number of nodes in the list

    void SLinkedList() {
        head = tail = null;
        size = 0;
    }

    void addFirst(Node n)  {...}

    Node removeFirst()  {...}

    ...
}
```



← all methods must update tail

Insert at tail

```
void addLast(Node n) {  
  
    //if list is empty the new element is head and tail  
    if (tail == null) {  
        n.setNext(null);  
        head = tail = n;  
    } else {  
        //the list is not empty: link tail to n and n becomes the new  
        //tail  
        tail.setNext(n);  
        n.setNext(null);  
        tail = n;  
    }  
    //increment size  
    size++  
}
```

- Special cases: list is empty
- Efficiency: Theta(1)

Remove at tail

- What we want: delete the last element and set the new tail
- Is that possible?

Remove at tail

- What we want: delete the last element and set the new tail
 - Is that possible?
 - Remove at tail
 - set the tail to the node BEFORE the tail
 - need the node before the tail: $O(\text{size})$
 - To remove an element from a list you need the node BEFORE it as well
- ```
remove(Node n) {
 //link n.before to n.next
}
```
- To remove a node efficiently need to keep track of previous node

# Doubly-linked lists



```
/** Node of a doubly linked list of integers */
public class DNode {
 protected int element; //element stored by a node
 protected DNode next, prev; // Pointers to next and previous nodes

 /** Constructor that creates a node with given fields */
 public DNode(int e, DNode p, DNode n) {
 element = e;
 prev = p;
 next = n;
 }
 /** Returns the element of this node */
 public int getElement() { return element; }
 /** Returns the previous node of this node */
 public DNode getPrev() { return prev; }
 /** Returns the next node of this node */
 public DNode getNext() { return next; }
 /** Sets the element of this node */
 public void setElement(Int newElem) { element = newElem; }
 /** Sets the previous node of this node */
 public void setPrev(DNode newPrev) { prev = newPrev; }
 /** Sets the next node of this node */
 public void setNext(DNode newNext) { next = newNext; }
}
```

# Doubly-linked lists

```
/** Doubly linked list with nodes of type DNode */
public class DList {

 protected int size; // number of elements
 protected DNode head, tail;

 void addFirst(Node n);
 void addLast(Node n);
 Node deleteFirst();
 Node deleteLast();
 void delete(Node n);
}
```

- Operations on doubly linked lists
  - `addFirst()`: O(1) time
  - `addLast()`: O(1) time
  - `deleteFirst()`: O(1) time
  - `deleteLast()`: O(1) time
  - `delete()`: O(1) time
  - `get(i)`: O(i) time

# Insert at head

```
void addFirst(Node n) {
 n.setNext(head);
 n.setprev(null);
 head.setPrev(n);
 head = n;
 size++;
}
```

Does this work?

# Insert at head

```
void addFirst(Node n) {

 n.setNext(head);

 n.setprev(null);

 head.setPrev(n);

 head = n;

 size++;

}
```

- Special cases?

- empty list: head is null; need to set tail too

- Efficiency ?
  - O(1)

```
void addFirst(Node n) {

 if (head==null) {

 /* this is the first
 element: set both head
 and tail to it */

 head = tail = n;

 n.setPrev(null);
 n.setNext(null);

 }

 else {

 n.setNext(head);
 n.setprev(null);

 head.setPrev(n);

 head = n;

 }

 size++;

}
```

# Insert at tail

```
void addLast(Node n) {
 tail.setNext(n);
 n.setprev(tail);
 n.setNect(null);
 tail = n;
 size++;
}
```

Does this work ?

# Insert at tail

```
void addLast(Node n) {

 tail.setNext(n);

 n.setprev(tail);

 n.setNect(null);

 tail = n;

 size++;

}
```

- Special cases?
  - empty list: tail is null; need to set head too
  - Efficiency: O(1)

```
void addLast(Node n) {

 if (tail == null) {

 head = tail = n;
 n.setPrev(null);
 n.setNext(null);

 }

 else {

 tail.setNext(n);
 n.setprev(tail);
 n.setNect(null);

 tail = n;

 }

 size++;

}
```

# Doubly-linked lists

- Class work: Sketch the following methods for doubly-linked lists, and analyze their efficiency.
  - `Node removeFirst()`
  - `Node removeLast()`
  - `void remove(Node n)`
  - `Node search(int k)`

# Sentinels

- Sentinels for singly-linked list: keep a dummy head
  - an empty list is one node: the dummy head
- Sentinels for doubly-linked lists
  - dummy head and dummy tail
- Why? elegant. Unifies special cases when head or tail are null

# DLLists with Sentinels

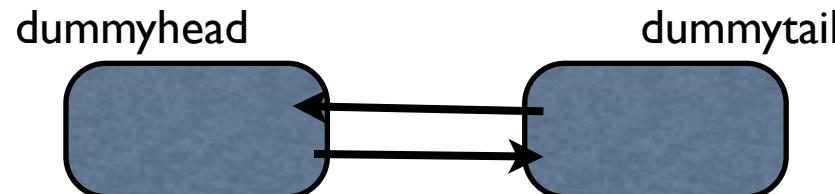
```
public class DList {

 protected int size; // number of elements
 protected DNode header, trailer; // sentinels

 /** Constructor that creates an empty list */
 public DList() {
 size = 0;
 header = new DNode(null, null, null); // create header
 trailer = new DNode(null, header, null); // create trailer
 // make header and trailer point to each other
 header.setNext(trailer);

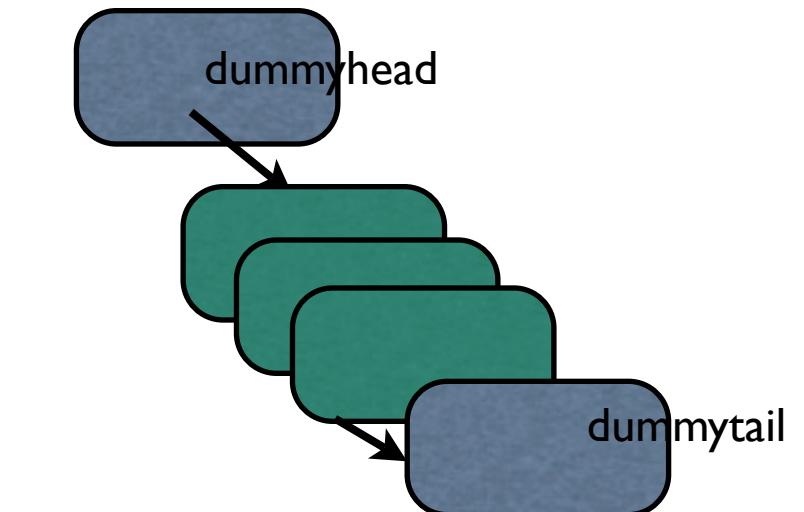
 }
}
```

- the empty list:
  - `size = 0`

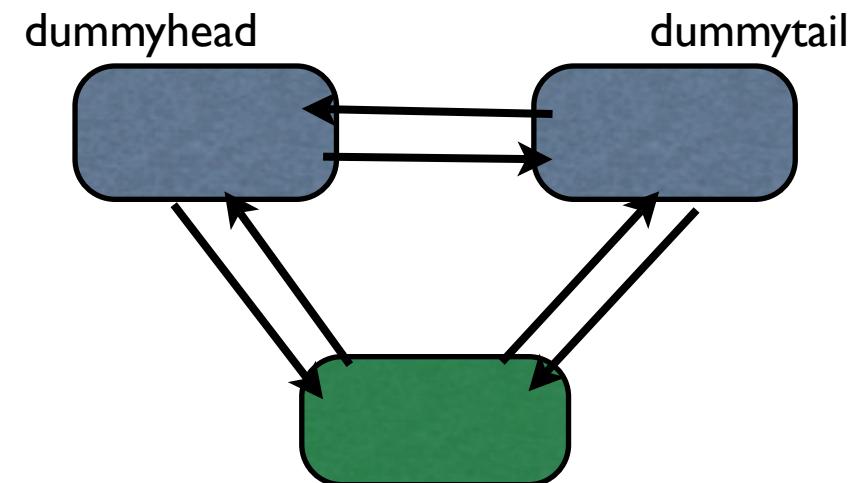


# DLLists with sentinels

```
insertFirst(Node n) {
 n.setNext(dummyHead.getNext());
 dummyHead.getNext().setPrev(n);
 dummyHead.setNext(n);
 n.setPrev(dummyhead);
 size++;
}
```



- Special cases: none
  - works for empty list



# Extensions

- Circular lists : make last node point to the first (instead of null)

```
class CircularList {

 SNode head;

 int size;

}
```

- Let's say we want to insert at head

```
insertAtHead(Node n) {

 n.setNext(head.getNext());

 head.setNext(n);

}
```

- If head is null?

```
if (head ==null) {

 n.setNext(n);

 head = n;

}
```

